



GD25R128E

DATASHEET



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1 FEATURES

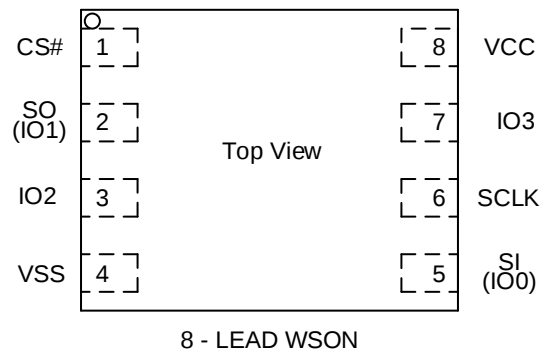
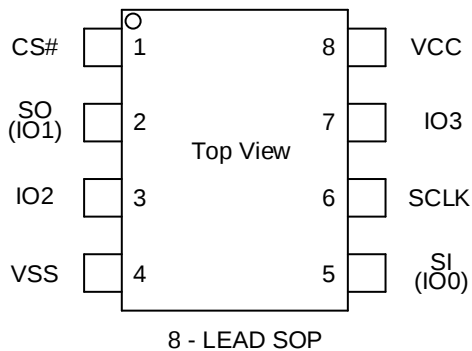
- ◆ 128M-bit Serial Flash
 - 16M-Byte
 - 256 Bytes per programmable page
- ◆ Standard, Dual, Quad SPI
 - Standard SPI: SCLK, CS#, SI, SO,
 - Dual SPI: SCLK, CS#, IO0, IO1,
 - Quad SPI: SCLK, CS#, IO0, IO1, IO2, IO3
- ◆ High Speed Clock Frequency
 - 104MHz for fast read with 30PF load
 - Dual I/O Data transfer up to 208Mbits/s
 - Quad I/O Data transfer up to 416Mbits/s
- ◆ Software Write Protection
 - Write protect all/portion of memory via software
 - Top/Bottom Block protection
- ◆ Endurance and Data Retention
 - Minimum 100,000 Program/Erase Cycles
 - 20-year data retention typical
- ◆ RPMC Function
 - Four 32-bit Monotonic Counters
 - Volatile HMAC Key Register
 - Non-volatile Rood Key Register
- ◆ Fast Program/Erase Speed
 - Page Program time: 0.5ms typical
 - Sector Erase time: 45ms typical
 - Block Erase time: 0.15s/0.25s typical
 - Chip Erase time: 50s typical
- ◆ Flexible Architecture
 - Uniform Sector of 4K-Byte
 - Uniform Block of 32/64K-Byte
- ◆ Low Power Consumption
 - 15 μ A typical standby current
 - 2 μ A typical deep power down current
- ◆ Advanced Security Features
 - 128-bit Unique ID for each device
 - Serial Flash Discoverable parameters (SFDP) register
 - 3x1024-Byte Security Registers With OTP Locks
- ◆ Single Power Supply Voltage
 - Full voltage range: 2.7-3.6V
- ◆ Package Information
 - SOP8 208mil
 - WSON8 (6x5mm)
 - WSON8 (8x6mm)



2 GENERAL DESCRIPTIONS

The GD25R128E (128M-bit) Serial flash supports the standard Serial Peripheral Interface (SPI), and the Dual/Quad SPI: Serial Clock, Chip Select, Serial Data I/O0 (SI), I/O1 (SO), I/O2, I/O3. The Dual I/O data is transferred with speed of 208Mbit/s, and the Quad I/O data is transferred with speed of 416Mbit/s.

CONNECTION DIAGRAM



PIN DESCRIPTION

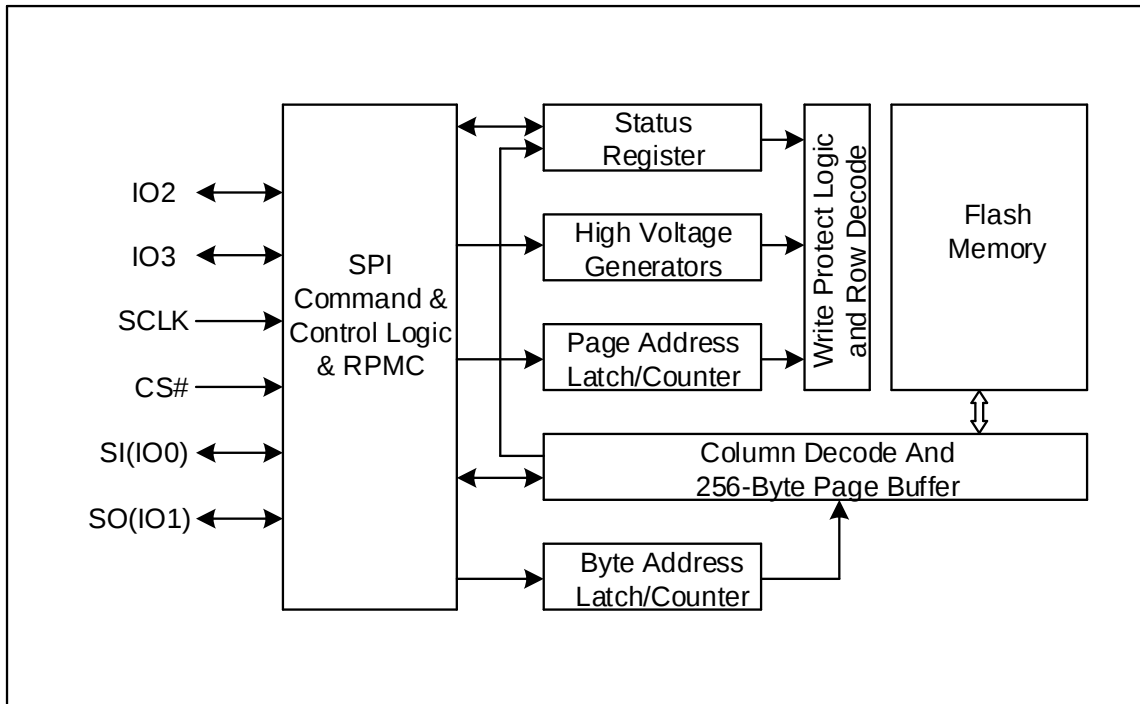
Table 1. Pin Description for SOP8/WSO8 Package

Pin No.	Pin Name	I/O	Description
1	CS#	I	Chip Select Input
2	SO (IO1)	I/O	Data Output (Data Input Output 1)
3	IO2	I/O	Data Input Output 2
4	VSS		Ground
5	SI (IO0)	I/O	Data Input (Data Input Output 0)
6	SCLK	I	Serial Clock Input
7	IO3	I/O	Data Input Output 3
8	VCC		Power Supply

Note: CS# must be driven high if chip is not selected. Please don't leave CS# floating any time after power is on.



BLOCK DIAGRAM





3 MEMORY ORGANIZATION

GD25R128E

Each device has	Each block has	Each sector has	Each page has	
16M	64/32K	4K	256	Bytes
64K	256/128	16	-	pages
4K	16/8	-	-	sectors
256/512	-	-	-	blocks

UNIFORM BLOCK SECTOR ARCHITECTURE

GD25R128E 64K Bytes Block Sector Architecture

Block	Sector	Address range	
255	4095	FFF000H	FFFFFFFH
			
	4080	FF0000H	FF0FFFFH
254	4079	FEF000H	FEFFFFFH
			
	4064	FE0000H	FE0FFFFH
.....			
			
			
.....			
			
			
2	47	02F000H	02FFFFFFH
			
	32	020000H	020FFFFH
1	31	01F000H	01FFFFFFH
			
	16	010000H	010FFFFH
0	15	00F000H	00FFFFFFH
			
	0	000000H	000FFFFH



4 DEVICE OPERATIONS

4.1 SPI Mode

Standard SPI

The GD25R128E features a serial peripheral interface on 4 signals bus: Serial Clock (SCLK), Chip Select (CS#), Serial Data Input (SI) and Serial Data Output (SO). Both SPI bus mode 0 and 3 are supported. Input data is latched on the rising edge of SCLK and data shifts out on the falling edge of SCLK.

Dual SPI

The GD25R128E supports Dual SPI operation when using the “Dual Output Fast Read” and “Dual I/O Fast Read” (3BH and BBH) commands. These commands allow data to be transferred to or from the device at twice the rate of the standard SPI. When using the Dual SPI commands, the SI and SO pins become bidirectional I/O pins: IO0 and IO1.

Quad SPI

The GD25R128E supports Quad SPI operation when using the “Quad Output Fast Read”, “Quad I/O Fast Read” (6BH, EBH) commands. These commands allow data to be transferred to or from the device at four times the rate of the standard SPI. When using the Quad SPI commands, the SI and SO pins become bidirectional I/O pins: IO0 and IO1, in addition to IO2 and IO3 pins. For GD25R128E, the QE bit is set to 1 as default and cannot be changed.



5 DATA PROTECTION

The GD25R128E provide the following data protection methods:

- ◆ Write Enable (WREN) command: The WREN command is set the Write Enable Latch bit (WEL). The WEL bit will return to reset by the following situation:
 - Power-Up / Software Reset (66H+99H)
 - Write Disable (WRDI)
 - Write Status Register (WRSR)
 - Page Program (PP)
 - Sector Erase (SE) / Block Erase (BE) / Chip Erase (CE)
- ◆ Software Protection Mode: The Block Protect bits (BP4-BP0) define the section of the memory array that can be read but not changed.
- ◆ Deep Power-Down Mode: In Deep Power-Down Mode, all commands are ignored except the Release from Deep Power-Down Mode command and Software Reset (66H+99H).
- ◆ Write Inhibit Voltage (VWI): Device would reset automatically when VCC is below a certain threshold VWI.

Table 2. GD25R128E Protected area size (CMP=0)

Status Register Content					Memory Content			
BP4	BP3	BP2	BP1	BP0	Blocks	Addresses	Density	Portion
X	X	0	0	0	NONE	NONE	NONE	NONE
0	0	0	0	1	252 to 255	FC0000H-FFFFFFH	256KB	Upper 1/64
0	0	0	1	0	248 to 255	F80000H-FFFFFFH	512KB	Upper 1/32
0	0	0	1	1	240 to 255	F00000H-FFFFFFH	1MB	Upper 1/16
0	0	1	0	0	224 to 255	E00000H-FFFFFFH	2MB	Upper 1/8
0	0	1	0	1	192 to 255	C00000H-FFFFFFH	4MB	Upper 1/4
0	0	1	1	0	128 to 255	800000H-FFFFFFH	8MB	Upper 1/2
0	1	0	0	1	0 to 3	000000H-03FFFFH	256KB	Lower 1/64
0	1	0	1	0	0 to 7	000000H-07FFFFH	512KB	Lower 1/32
0	1	0	1	1	0 to 15	000000H-0FFFFFH	1MB	Lower 1/16
0	1	1	0	0	0 to 31	000000H-1FFFFFH	2MB	Lower 1/8
0	1	1	0	1	0 to 63	000000H-3FFFFFH	4MB	Lower 1/4
0	1	1	1	0	0 to 127	000000H-7FFFFFH	8MB	Lower 1/2
X	X	1	1	1	0 to 255	000000H-FFFFFFH	16MB	ALL
1	0	0	0	1	255	FFF000H-FFFFFFH	4KB	Top Block
1	0	0	1	0	255	FFE000H-FFFFFFH	8KB	Top Block
1	0	0	1	1	255	FFC000H-FFFFFFH	16KB	Top Block
1	0	1	0	X	255	FF8000H-FFFFFFH	32KB	Top Block
1	0	1	1	0	255	FF8000H-FFFFFFH	32KB	Top Block
1	1	0	0	1	0	000000H-000FFFFH	4KB	Bottom Block
1	1	0	1	0	0	000000H-001FFFFH	8KB	Bottom Block
1	1	0	1	1	0	000000H-003FFFFH	16KB	Bottom Block
1	1	1	0	X	0	000000H-007FFFFH	32KB	Bottom Block
1	1	1	1	0	0	000000H-007FFFFH	32KB	Bottom Block



Table 3. GD25R128E Protected area size (CMP=1)

Status Register Content					Memory Content			
BP4	BP3	BP2	BP1	BP0	Blocks	Addresses	Density	Portion
X	X	0	0	0	0 to 255	000000H-FFFFFFH	ALL	ALL
0	0	0	0	1	0 to 251	000000H-FBFFFFH	16128KB	Lower 63/64
0	0	0	1	0	0 to 247	000000H-F7FFFFH	15872KB	Lower 31/32
0	0	0	1	1	0 to 239	000000H-EFFFFFH	15MB	Lower 15/16
0	0	1	0	0	0 to 223	000000H-DFFFFFH	14MB	Lower 7/8
0	0	1	0	1	0 to 191	000000H-BFFFFFH	12MB	Lower 3/4
0	0	1	1	0	0 to 127	000000H-7FFFFFH	8MB	Lower 1/2
0	1	0	0	1	4 to 255	040000H-FFFFFFH	16128KB	Upper 63/64
0	1	0	1	0	8 to 255	080000H-FFFFFFH	15872KB	Upper 31/32
0	1	0	1	1	16 to 255	100000H-FFFFFFH	15MB	Upper 15/16
0	1	1	0	0	32 to 255	200000H-FFFFFFH	14MB	Upper 7/8
0	1	1	0	1	64 to 255	400000H-FFFFFFH	12MB	Upper 3/4
0	1	1	1	0	128 to 255	800000H-FFFFFFH	8MB	Upper 1/2
X	X	1	1	1	NONE	NONE	NONE	NONE
1	0	0	0	1	0 to 255	000000H-FFEFFFFH	16380KB	L-4095/4096
1	0	0	1	0	0 to 255	000000H-FFDFFFFH	16376KB	L-2047/2048
1	0	0	1	1	0 to 255	000000H-FFBFFFFH	16368KB	L-1023/1024
1	0	1	0	X	0 to 255	000000H-FF7FFFFH	16352KB	L-511/512
1	0	1	1	0	0 to 255	000000H-FF7FFFFH	16352KB	L-511/512
1	1	0	0	1	0 to 255	001000H-FFFFFFH	16380KB	U-4095/4096
1	1	0	1	0	0 to 255	002000H-FFFFFFH	16376KB	U-2047/2048
1	1	0	1	1	0 to 255	004000H-FFFFFFH	16368KB	U-1023/1024
1	1	1	0	X	0 to 255	008000H-FFFFFFH	16352KB	U-511/512
1	1	1	1	0	0 to 255	008000H-FFFFFFH	16352KB	U-511/512



6 STATUS REGISTER

Table 4. Status Register-SR No.1

No.	Name	Description	Note
S7	SRP0	Status Register Protection Bit	Non-volatile writable
S6	BP4	Block Protect Bit	Non-volatile writable
S5	BP3	Block Protect Bit	Non-volatile writable
S4	BP2	Block Protect Bit	Non-volatile writable
S3	BP1	Block Protect Bit	Non-volatile writable
S2	BP0	Block Protect Bit	Non-volatile writable
S1	WEL	Write Enable Latch	Volatile, read only
S0	WIP	Erase/Write In Progress	Volatile, read only

Table 5. Status Register-SR No.2

No.	Name	Description	Note
S15	SUS1	Erase Suspend Bit	Volatile, read only
S14	CMP	Complement Protect Bit	Non-volatile writable
S13	LB3	Security Register Lock Bit	Non-volatile writable (OTP)
S12	LB2	Security Register Lock Bit	Non-volatile writable (OTP)
S11	LB1	Security Register Lock Bit	Non-volatile writable (OTP)
S10	SUS2	Program Suspend Bit	Volatile, read only
S9	QE	Quad Enable Bit	QE = 1 permanently
S8	SRP1	Status Register Protection Bit	Non-volatile writable

Table 6. Status Register-SR No.3

No.	Name	Description	Note
S23	Reserved	Reserved	Reserved
S22	DRV1	Output Driver Strength Bit	Non-volatile writable
S21	DRV0	Output Driver Strength Bit	Non-volatile writable
S20	Reserved	Reserved	Reserved
S19	Reserved	Reserved	Reserved
S18	Reserved	Reserved	Reserved
S17	Reserved	Reserved	Reserved
S16	DC	Dummy Configuration Bit	Non-volatile writable

The status and control bits of the Status Register are as follows:

WIP bit

The Write in Progress (WIP) bit indicates whether the memory is busy in program/erase/write status register progress. When WIP bit sets to 1, means the device is busy in program/erase/write status register progress, when WIP bit sets 0, means the device is not in program/erase/write status register progress.



WEL bit

The Write Enable Latch (WEL) bit indicates the status of the internal Write Enable Latch. When set to 1 the internal Write Enable Latch is set, when set to 0 the internal Write Enable Latch is reset and no Write Status Register, Program or Erase command is accepted.

BP4, BP3, BP2, BP1, BP0 bits

The Block Protect (BP4, BP3, BP2, BP1, and BP0) bits are non-volatile. They define the size of the area to be software protected against Program and Erase commands. These bits are written with the Write Status Register (WRSR) command. When the Block Protect (BP4, BP3, BP2, BP1, BP0) bits are set to 1, the relevant memory area (as defined in Table 2&3) becomes protected against Page Program (PP), Sector Erase (SE) and Block Erase (BE) commands. The Chip Erase (CE) command is executed, if the Block Protect (BP2, BP1, and BP0) bits are 0 and CMP=0 or the Block Protect (BP2, BP1, and BP0) bits are 1 and CMP=1.

SRP1, SRP0 bits

The Status Register Protect (SRP1 and SRP0) bits are non-volatile Read/Write bits in the status register. The SRP bits control the method of write protection: software protection, power supply lock-down or one time programmable protection.

SRP1	SRP0	Status Register	Description
0	0	Software Protected	The Status Register can be written to after a Write Enable command, WEL=1.(Default)
1	0	Power Supply Lock-Down ⁽¹⁾⁽²⁾	Status Register is protected and cannot be written to again until the next Power-Down, Power-Up cycle.
1	1	One Time Program ⁽²⁾	Status Register is permanently protected and cannot be written to.

NOTE:

1. When SRP1, SRP0= (1, 0), a Power-Down, Power-Up cycle will change SRP1, SRP0 to (0, 0) state.
2. This feature is available on special order. Please contact GigaDevice for details.

QE bit

The Quad Enable (QE) bit is a non-volatile bit in the Status Register that allows Quad operation. The default value of QE bit is 1 and it cannot be changed, so that the IO2 and IO3 pins are enabled all the time.

LB3, LB2, LB1 bits

The LB3, LB2 and LB1 bits are non-volatile One Time Program (OTP) bits in Status Register (S13, S12 and S11) that provide the write protect control and status to the Security Registers. The default state of LB3, LB2 and LB1 bits are 0, the security registers are unlocked. The LB3, LB2 and LB1 bits can be set to 1 individually using the Write Register instruction. The LB3, LB2 and LB1 bits are One Time Programmable, once they are set to 1, the Security Registers will become read-only permanently.

CMP bit

The CMP bit is a non-volatile Read/Write bit in the Status Register (S14). It is used in conjunction with the BP4-BP0 bits to provide more flexibility for the array protection. Please see the Status registers Memory Protection table for details. The default setting is CMP=0.



SUS1, SUS2 bits

The SUS1 and SUS2 bits are read only bits in the status register (S15 and S10) that are set to 1 after executing an Erase/Program Suspend (75H) command (The Erase Suspend will set the SUS1 bit to 1, and the Program Suspend will set the SUS2 bit to 1). The SUS1 and SUS2 bits are cleared to 0 by Erase/Program Resume (7AH) command, software reset (66H+99H) command, as well as a power-down, power-up cycle.

DC bit

The Dummy Configuration (DC) bit is non-volatile, which selects the number of dummy cycles between the end of address and the start of read data output. Dummy cycles provide additional latency that is needed to complete the initial read access of the flash array before data can be returned to the host system. Some read commands require additional dummy cycles as the SCLK frequency increases.

The following dummy cycle tables provide different dummy cycle settings that are configured.

Command	DC bit	Numbers of Dummy Cycles	Freq.(MHz)
BBH	0 (default)	4	104
	1	8	104
EBH	0 (default)	6	104
	1	10	104

DRV1, DRV0 bits

The DRV1 and DRV0 bits are used to determine the output driver strength for the Read operations.

Table 7. Driver Strength for Read Operations

DRV1, DRV0	Driver Strength
00	100%
01	75% (default)
10	50%
11	25%

Reserved bit

It is recommended to set the value of the reserved bit as "0".



7 COMMAND DESCRIPTIONS

All commands, addresses and data are shifted in and out of the device, beginning with the most significant bit on the first rising edge of SCLK after CS# is driven low. Then, the one-byte command code must be shifted in to the device, with most significant bit first on SI, and each bit is latched on the rising edges of SCLK.

Every command sequence starts with a one-byte command code. Depending on the command, this might be followed by address bytes, or by data bytes, or by both or none. CS# must be driven high after the last bit of the command sequence has been completed. For the command of Read, Fast Read, Read Status Register or Release from Deep Power-Down, and Read Device ID, the shifted-in command sequence is followed by a data-out sequence. All read instruction can be completed after any bit of the data-out sequence is being shifted out, and then CS# must be driven high to return to deselected status.

For the command of Page Program, Sector Erase, Block Erase, Chip Erase, Write Status Register, Write Enable, Write Disable or Deep Power-Down command, CS# must be driven high exactly at a byte boundary, otherwise the command is rejected, and is not executed. That is CS# must be driven high when the number of clock pulses after CS# being driven low is an exact multiple of eight. For Page Program, if at any time the input byte is not a full byte, nothing will happen and WEL will not be reset.

Table 8. Commands

Command Name	Byte 1	Byte 2	Byte 3	Byte 4	Byte 5	Byte 6	Byte 7	Byte 8	Byte 9
Write Enable	06H								
Write Disable	04H								
Read Status Register-1	05H	(S7-S0)	(cont.)						
Read Status Register-2	35H	(S15-S8)	(cont.)						
Read Status Register-3	15H	(S23-S16)	(cont.)						
Write Status Register-1	01H	S7-S0							
Write Status Register-2	31H	S15-S8							
Write Status Register-3	11H	S23-S16							
Volatile SR write Enable	50H								
Read Data	03H	A23-A16	A15-A8	A7-A0	(D7-D0)	(cont.)			
Fast Read	0BH	A23-A16	A15-A8	A7-A0	dummy	(D7-D0)	(cont.)		
Dual Output Fast Read	3BH	A23-A16	A15-A8	A7-A0	dummy	(D7-D0) ⁽¹⁾	(cont.)		
Quad Output Fast Read	6BH	A23-A16	A15-A8	A7-A0	dummy	(D7-D0) ⁽²⁾	(cont.)		
Dual I/O Fast Read	BBH	A23-A16 ⁽³⁾	A15-A8 ⁽³⁾	A7-A0 ⁽³⁾	M7-M0 ⁽⁴⁾	(D7-D0) ⁽¹⁾	(cont.)		
Quad I/O Fast Read	EBH	A23-A16 ⁽⁵⁾	A15-A8 ⁽⁵⁾	A7-A0 ⁽⁵⁾	M7-M0 ⁽⁶⁾	dummy	dummy	(D7-D0) ⁽²⁾	(cont.)
Set Burst with Wrap	77H	dummy ⁽⁷⁾	dummy ⁽⁷⁾	dummy ⁽⁷⁾	W7-W0 ⁽⁷⁾				
Page Program	02H	A23-A16	A15-A8	A7-A0	D7-D0	Next Byte			



Uniform Sector GigaDevice Dual and Quad Serial Flash

GD25R128E

Quad Page Program	32H	A23-A16	A15-A8	A7-A0	D7-D0 ⁽⁸⁾	Next Byte			
Sector Erase	20H	A23-A16	A15-A8	A7-A0					
Block Erase (32K)	52H	A23-A16	A15-A8	A7-A0					
Block Erase (64K)	D8H	A23-A16	A15-A8	A7-A0					
Chip Erase	C7/60H								
Read Manufacturer/ Device ID	90H	00H	00H	00H	(MID7- MID0)	(ID7-ID0)	(cont.)		
Read Identification	9FH	(MID7- MID0)	(ID15-ID8)	(ID7-ID0)	(cont.)				
Read Unique ID	4BH	00H	00H	00H	dummy	(UID7- UID0)	(cont.)		
Erase Security Registers ⁽⁹⁾	44H	A23-A16	A15-A8	A7-A0					
Program Security Registers ⁽⁹⁾	42H	A23-A16	A15-A8	A7-A0	D7-D0	Next Byte			
Read Security Registers ⁽⁹⁾	48H	A23-A16	A15-A8	A7-A0	dummy	(D7-D0)	(cont.)		
Enable Reset	66H								
Reset	99H								
Program/Erase Suspend	75H								
Program/Erase Resume	7AH								
Deep Power-Down	B9H								
Release From Deep Power-Down	ABH								
Release From Deep Power-Down and Read Device ID	ABH	dummy	dummy	dummy	(ID7-ID0)	(cont.)			
Read Serial Flash Discoverable Parameter	5AH	A23-A16	A15-A8	A7-A0	dummy	(D7-D0)	(cont.)		

Note:

1. Dual Output data

IO0 = (D6, D4, D2, D0)

IO1 = (D7, D5, D3, D1)

2. Quad Output Data

IO0 = (D4, D0, ...)

IO1 = (D5, D1, ...)

IO2 = (D6, D2, ...)

IO3 = (D7, D3, ...)

3. Dual Input Address



IO0 = A22, A20, A18, A16, A14, A12, A10, A8 A6, A4, A2, A0

IO1 = A23, A21, A19, A17, A15, A13, A11, A9 A7, A5, A3, A1

4. Dual Input Mode bit

IO0 = M6, M4, M2, M0

IO1 = M7, M5, M3, M1

5. Quad Input Address

IO0 = A20, A16, A12, A8, A4, A0

IO1 = A21, A17, A13, A9, A5, A1

IO2 = A22, A18, A14, A10, A6, A2

IO3 = A23, A19, A15, A11, A7, A3

6. Quad Input Mode bit

IO0 = M4, M0

IO1 = M5, M1

IO2 = M6, M2

IO3 = M7, M3

7. Dummy bits and Wrap Bits

IO0 = (x, x, x, x, x, x, W4, x)

IO1 = (x, x, x, x, x, x, W5, x)

IO2 = (x, x, x, x, x, x, W6, x)

IO3 = (x, x, x, x, x, x, x, x)

8. Quad Output Data

IO0 = D4, D0, ...

IO1 = D5, D1, ...

IO2 = D6, D2, ...

IO3 = D7, D3, ...

9. Security Registers Address

Security Register1: A23-A16=00H, A15-A12=1H, A11-A10 = 00b, A9-A0= Byte Address;

Security Register2: A23-A16=00H, A15-A12=2H, A11-A10 = 00b, A9-A0= Byte Address;

Security Register3: A23-A16=00H, A15-A12=3H, A11-A10 = 00b, A9-A0= Byte Address;

TABLE OF ID DEFINITIONS

GD25R128E

Operation Code	MID7-MID0	ID15-ID8	ID7-ID0
9FH	C8	40	18
90H	C8		17
ABH			17

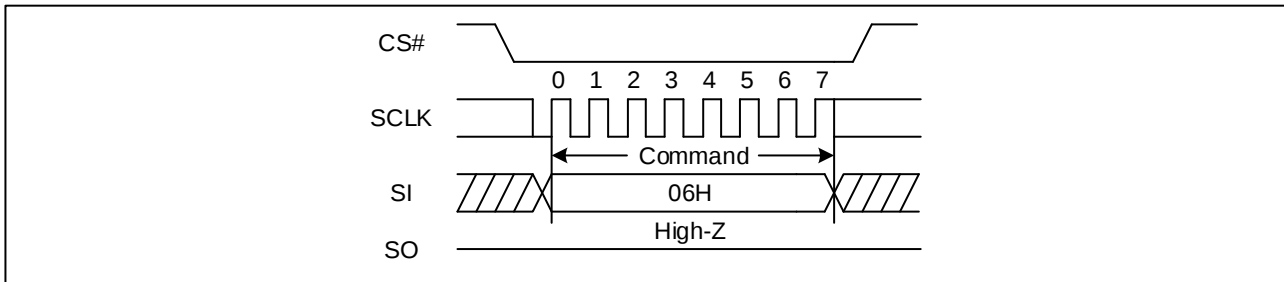


7.1 Write Enable (WREN) (06H)

The Write Enable (WREN) command is for setting the Write Enable Latch (WEL) bit. The Write Enable Latch (WEL) bit must be set prior to every Page Program (PP), Quad Page Program (QPP), Sector Erase (SE), Block Erase (BE), Chip Erase (CE), Write Status Register (WRSR) and Erase/Program Security Registers command.

The Write Enable (WREN) command sequence: CS# goes low → sending the Write Enable command → CS# goes high.

Figure 1. Write Enable Sequence Diagram



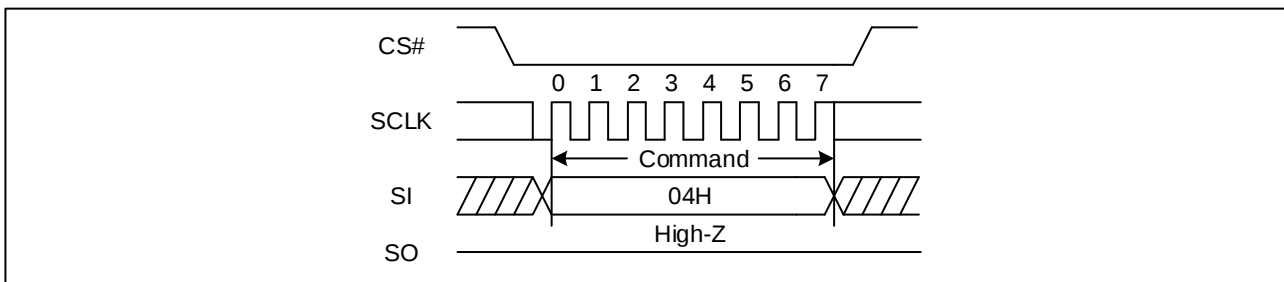
7.2 Write Disable (WRDI) (04H)

The Write Disable command is for resetting the Write Enable Latch (WEL) bit. The Write Enable Latch (WEL) bit may be set to 0 by issuing the Write Disable (WRDI) command to disable Page Program (PP), Quad Page Program (QPP), Sector Erase (SE), Block Erase (BE), Chip Erase (CE), Write Status Register (WRSR), that require WEL be set to 1 for execution. The WRDI command can be used by the user to protect memory areas against inadvertent writes that can possibly corrupt the contents of the memory. The WRDI command is ignored during an embedded operation while WIP bit =1.

The WEL bit is reset by following condition: Write Disable command (WRDI), Power-up, and upon completion of the Write Status Register, Page Program, Sector Erase, Block Erase and Chip Erase commands.

The Write Disable command sequence: CS# goes low → Sending the Write Disable command → CS# goes high.

Figure 2. Write Disable Sequence Diagram

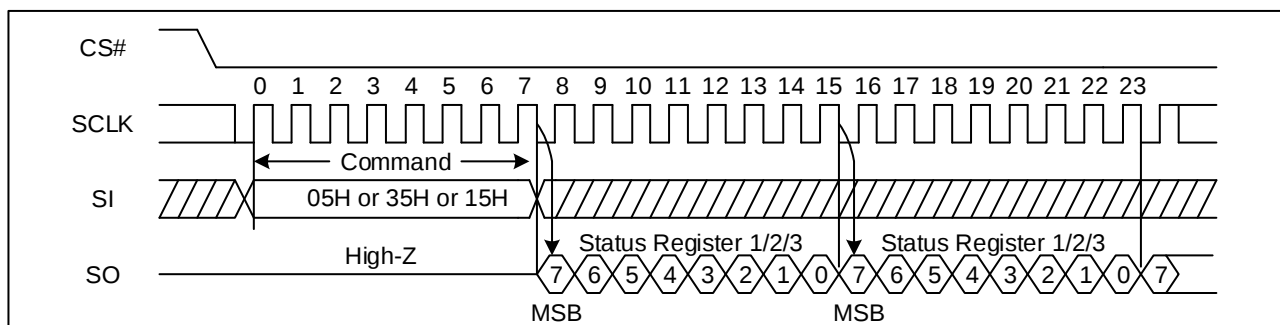


7.3 Read Status Register (RDSR) (05H or 35H or 15H)

The Read Status Register (RDSR) command is for reading the Status Register. The Status Register may be read at any time, even while a Program, Erase or Write Status Register cycle is in progress. When one of these cycles is in progress, it is recommended to check the Write in Progress (WIP) bit before sending a new command to the device. It is also possible to read the Status Register continuously. For command code of "05H" / "35H" / "15H", the SO will output Status Register bits S7~S0 / S15~S8 / S23~S16.



Figure 3. Read Status Register Sequence Diagram



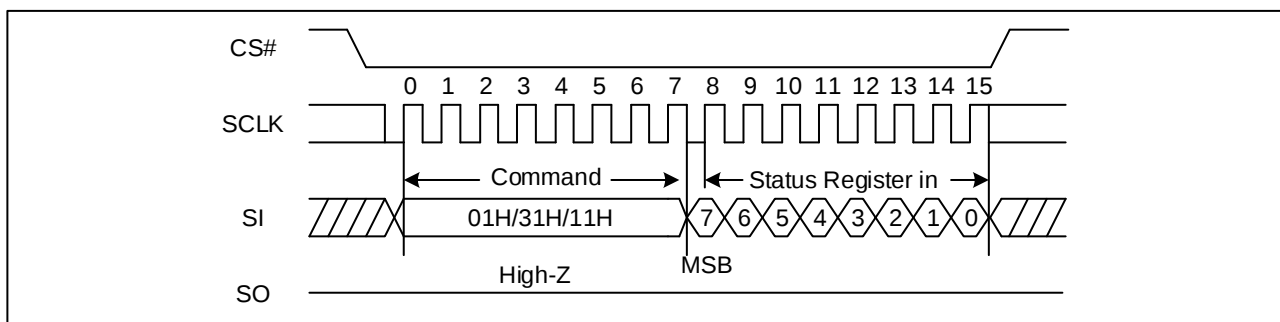
7.4 Write Status Register (WRSR) (01H or 31H or 11H)

The Write Status Register (WRSR) command allows new values to be written to the Status Register. Before it can be accepted, a Write Enable (WREN) command must previously have been executed. After the Write Enable (WREN) command has been decoded and executed, the device sets the Write Enable Latch (WEL).

The Write Status Register (WRSR) command has no effect on S15, S10, S9, S1 and S0 of the Status Register. For command code of "01H" / "31H" / "11H", the Status Register bits S7~S0 / S15~S8 / S23~S16 would be written. CS# must be driven high after the eighth bit of the data byte has been latched in. Otherwise, the Write Status Register (WRSR) command is not executed. As soon as CS# is driven high, the self-timed Write Status Register cycle (whose duration is t_w) is initiated. While the Write Status Register cycle is in progress, the Status Register may still be read to check the value of the Write In Progress (WIP) bit. The Write In Progress (WIP) bit is 1 during the self-timed Write Status Register cycle, and is 0 when it is completed. When the cycle is completed, the Write Enable Latch (WEL) is reset.

The Write Status Register (WRSR) command allows the user to change the values of the Block Protect (BP4, BP3, BP2, BP1, and BP0) bits, to define the size of the area that is to be treated as read-only.

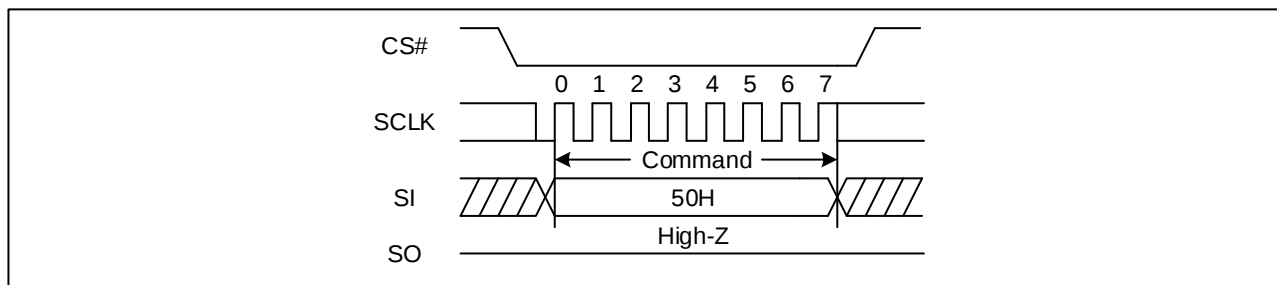
Figure 4. Write Status Register Sequence Diagram



7.5 Write Enable for Volatile Status Register (50H)

The non-volatile Status Register bits can also be written to as volatile bits. This gives more flexibility to change the system configuration and memory protection schemes quickly without waiting for the typical non-volatile bit write cycles or affecting the endurance of the Status Register non-volatile bits. The Write Enable for Volatile Status Register command must be issued prior to a Write Status Register command and any other commands can't be inserted between them. Otherwise, Write Enable for Volatile Status Register will be cleared. The Write Enable for Volatile Status Register command will not set the Write Enable Latch bit, it is only valid for the Write Status Register command to change the volatile Status Register bit values.

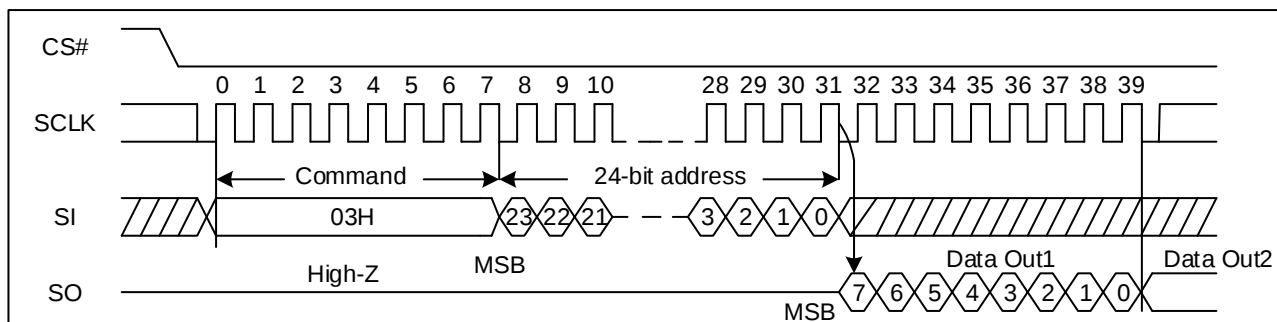
Figure 5. Write Enable for Volatile Status Register Sequence Diagram



7.6 Read Data Bytes (READ) (03H)

The Read Data Bytes (READ) command is followed by a 3-byte address (A23-A0), and each bit is latched-in on the rising edge of SCLK. Then the memory content, at that address, is shifted out on SO, and each bit is shifted out, at a Max frequency f_R , on the falling edge of SCLK. The first byte addressed can be at any location. The address is automatically incremented to the next higher address after each byte of data is shifted out. The whole memory can, therefore, be read with a single Read Data Bytes (READ) command. Any Read Data Bytes (READ) command, while an Erase, Program or Write cycle is in progress, is rejected without having any effects on the cycle that is in progress.

Figure 6. Read Data Bytes Sequence Diagram

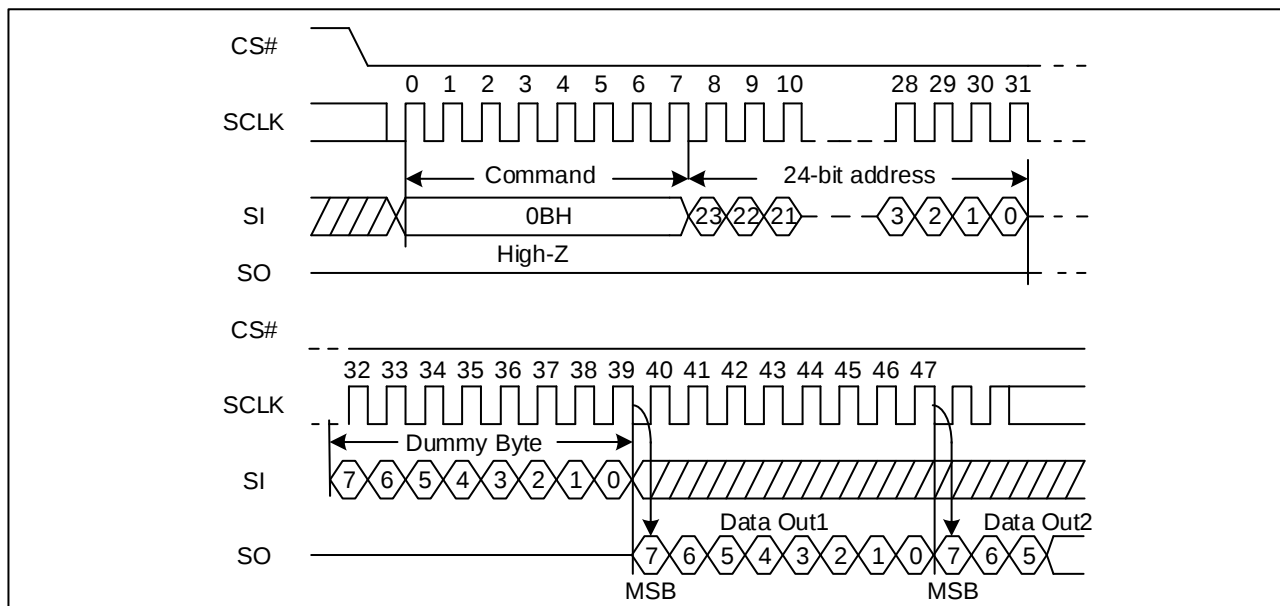


7.7 Read Data Bytes at Higher Speed (Fast Read) (0BH)

The Read Data Bytes at Higher Speed (Fast Read) command is for quickly reading data out. It is followed by a 3-byte address (A23-A0) and a dummy byte, and each bit is latched-in on the rising edge of SCLK. Then the memory content, at that address, is shifted out on SO, and each bit is shifted out, at a Max frequency f_c , on the falling edge of SCLK. The first byte address can be at any location. The address is automatically incremented to the next higher address after each byte of data is shifted out.



Figure 7. Read Data Bytes at Higher Speed Sequence Diagram

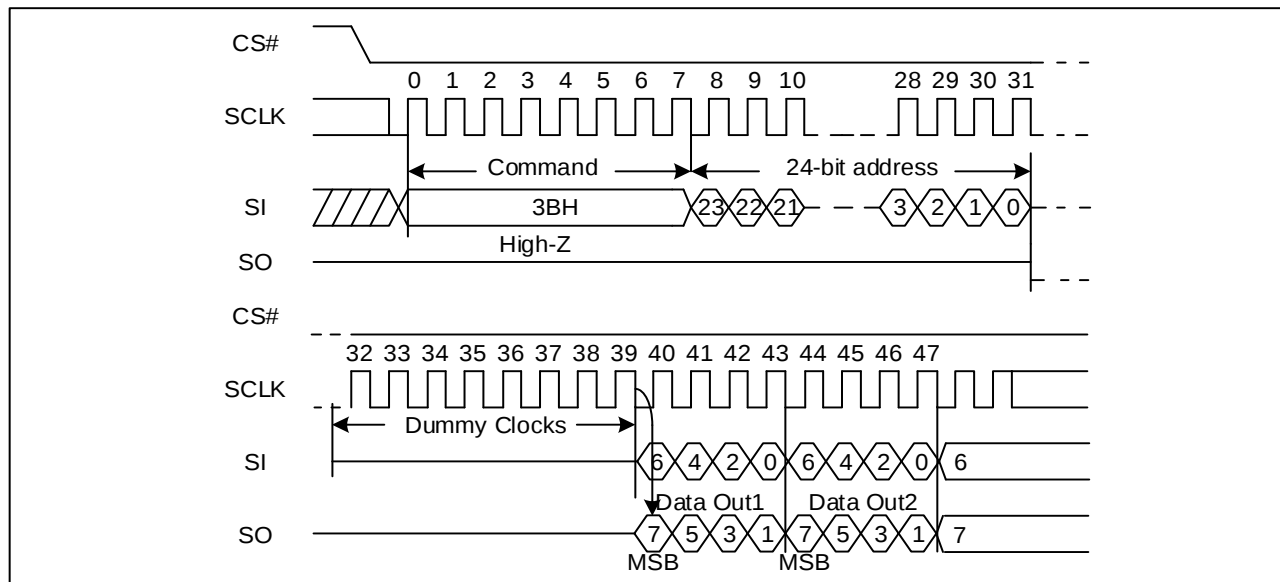


7.8 Dual Output Fast Read (3BH)

The Dual Output Fast Read command is followed by 3-byte address (A23-A0) and a dummy byte, and each bit is latched in on the rising edge of SCLK, then the memory contents are shifted out 2-bit per clock cycle from SI and SO.

The first byte addressed can be at any location. The address is automatically incremented to the next higher address after each byte of data is shifted out.

Figure 8. Dual Output Fast Read Sequence Diagram

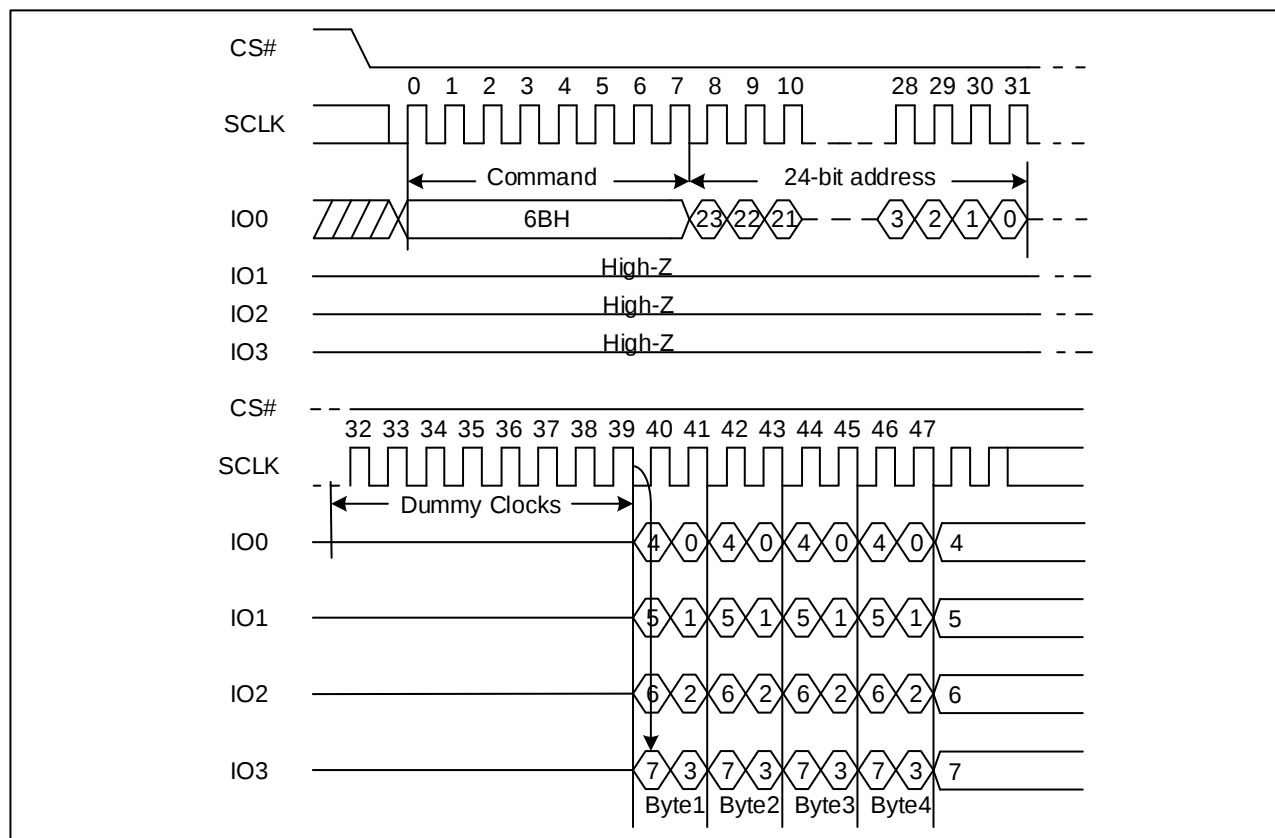




7.9 Quad Output Fast Read (6BH)

The Quad Output Fast Read command is followed by 3-byte address (A23-A0) and a dummy byte, and each bit is latched in on the rising edge of SCLK, then the memory contents are shifted out 4-bit per clock cycle from IO3, IO2, IO1 and IO0. The first byte addressed can be at any location. The address is automatically incremented to the next higher address after each byte of data is shifted out.

Figure 9. Quad Output Fast Read Sequence Diagram



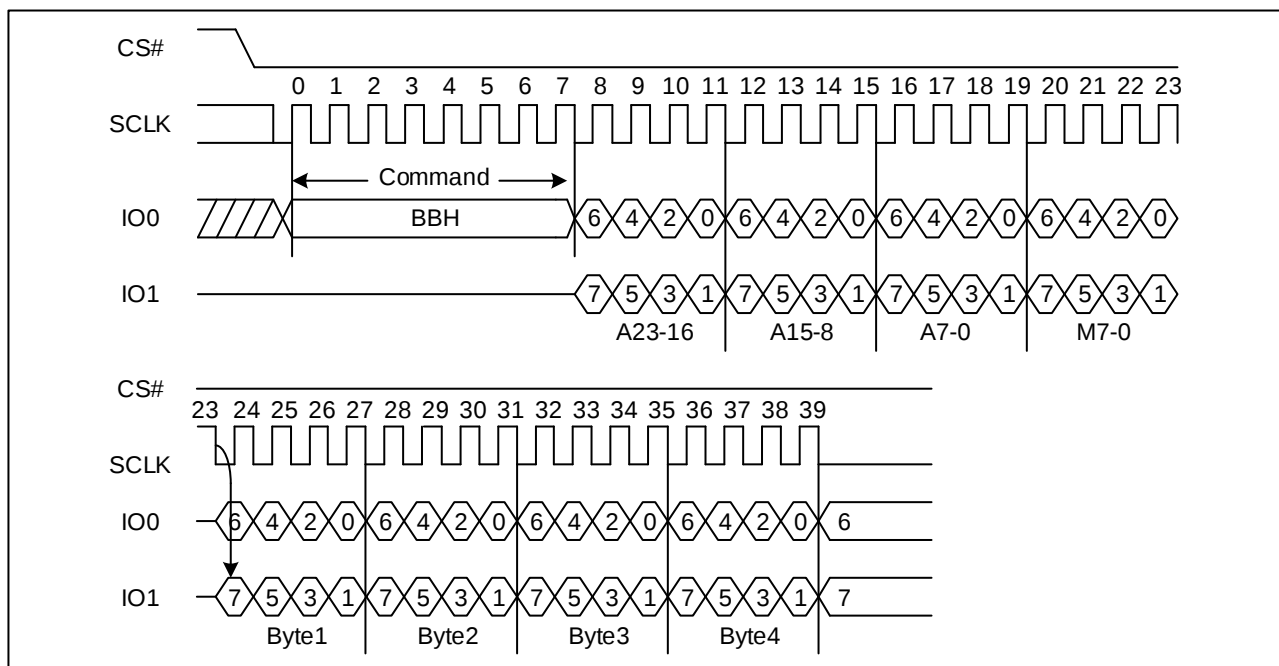
7.10 Dual I/O Fast Read (BBH)

The Dual I/O Fast Read command is similar to the Dual Output Fast Read command but with the capability to input the 3-byte address (A23-0) and a "Continuous Read Mode" byte 2-bit per clock by SI and SO, and each bit is latched in on the rising edge of SCLK, then the memory contents are shifted out 2-bit per clock cycle from SI and SO. The first byte addressed can be at any location. The address is automatically incremented to the next higher address after each byte of data is shifted out.

"Continuous Read Mode" bits must NOT be set as (M5-4) = (1, 0)



Figure 10. Dual I/O Fast Read Sequence Diagram ((M5-4) ≠ (1, 0))

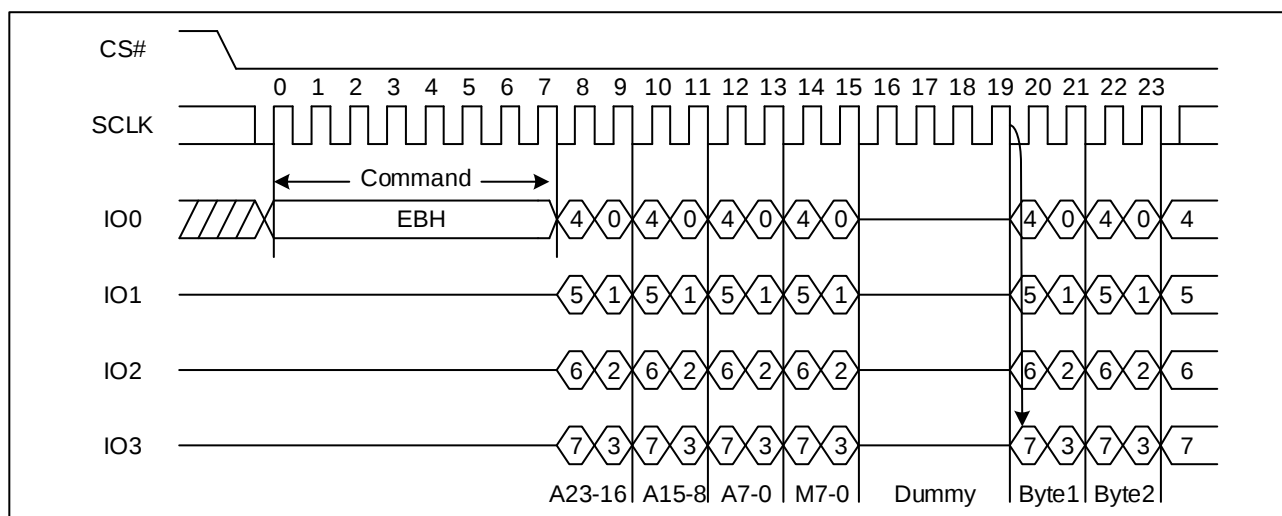


7.11 Quad I/O Fast Read (EBH)

The Quad I/O Fast Read command is similar to the Dual I/O Fast Read command but with the capability to input the 3-byte address (A23-0) and a “Continuous Read Mode” byte and 4-dummy clock 4-bit per clock by IO0, IO1, IO2, IO3, and each bit is latched in on the rising edge of SCLK, then the memory contents are shifted out 4-bit per clock cycle from IO0, IO1, IO2, IO3. The first byte addressed can be at any location. The address is automatically incremented to the next higher address after each byte of data is shifted out.

“Continuous Read Mode” bits must NOT be set as (M5-4) = (1, 0)

Figure 11. Quad I/O Fast Read Sequence Diagram ((M5-4) ≠ (1, 0))



Quad I/O Fast Read with “8/16/32/64-Byte Wrap Around” in Standard SPI mode

The Quad I/O Fast Read command can be used to access a specific portion within a page by issuing “Set Burst with Wrap” (77H) commands prior to EBH. The “Set Burst with Wrap” (77H) command can either enable or disable the “Wrap Around”



feature for the following EBH commands. When “Wrap Around” is enabled, the data being accessed can be limited to either an 8/16/32/64-byte section of a 256-byte page. The output data starts at the initial address specified in the command, once it reaches the ending boundary of the 8/16/32/64-byte section, the output will wrap around the beginning boundary automatically until CS# is pulled high to terminate the command.

The Burst with Wrap feature allows applications that use cache to quickly fetch a critical address and then fill the cache afterwards within a fixed length (8/16/32/64-byte) of data without issuing multiple read commands. The “Set Burst with Wrap” command allows three “Wrap Bits” W6-W4 to be set. The W4 bit is used to enable or disable the “Wrap Around” operation while W6-W5 is used to specify the length of the wrap around section within a page.

7.12 Set Burst with Wrap (77H)

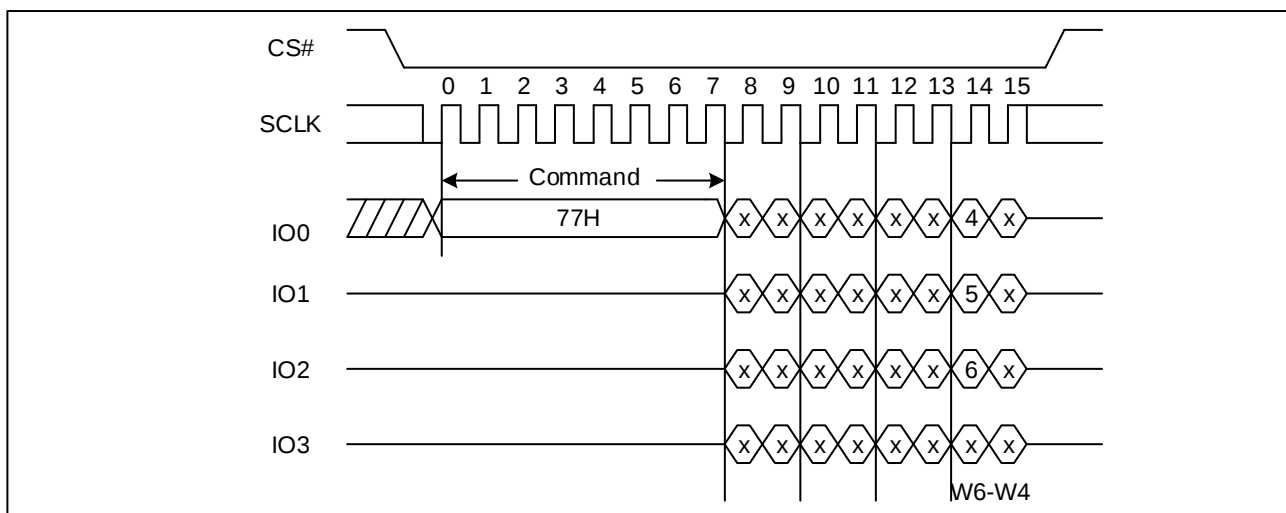
The Set Burst with Wrap command is used in conjunction with “Quad I/O Fast Read” command to access a fixed length of 8/16/32/64-byte section within a 256-byte page, in standard SPI mode.

The Set Burst with Wrap command sequence: CS# goes low → Send Set Burst with Wrap command → Send 24 dummy bits → Send 8 bits “Wrap bits” → CS# goes high.

W6,W5	W4=0		W4=1 (default)	
	Wrap Around	Wrap Length	Wrap Around	Wrap Length
0, 0	Yes	8-byte	No	N/A
0, 1	Yes	16-byte	No	N/A
1, 0	Yes	32-byte	No	N/A
1, 1	Yes	64-byte	No	N/A

If the W6-W4 bits are set by the Set Burst with Wrap command, all the following “Quad I/O Fast Read” command will use the W6-W4 setting to access the 8/16/32/64-byte section within any page. To exit the “Wrap Around” function and return to normal read operation, another Set Burst with Wrap command should be issued to set W4=1.

Figure 12. Set Burst with Wrap Sequence Diagram



7.13 Page Program (PP) (02H)

The Page Program (PP) command is for programming the memory. A Write Enable (WREN) command must previously have been executed to set the Write Enable Latch (WEL) bit before sending the Page Program command.

The Page Program (PP) command is entered by driving CS# Low, followed by the command code, three address bytes and at least one data byte on SI. If the 8 least significant address bits (A7-A0) are not all zero, all transmitted data that goes

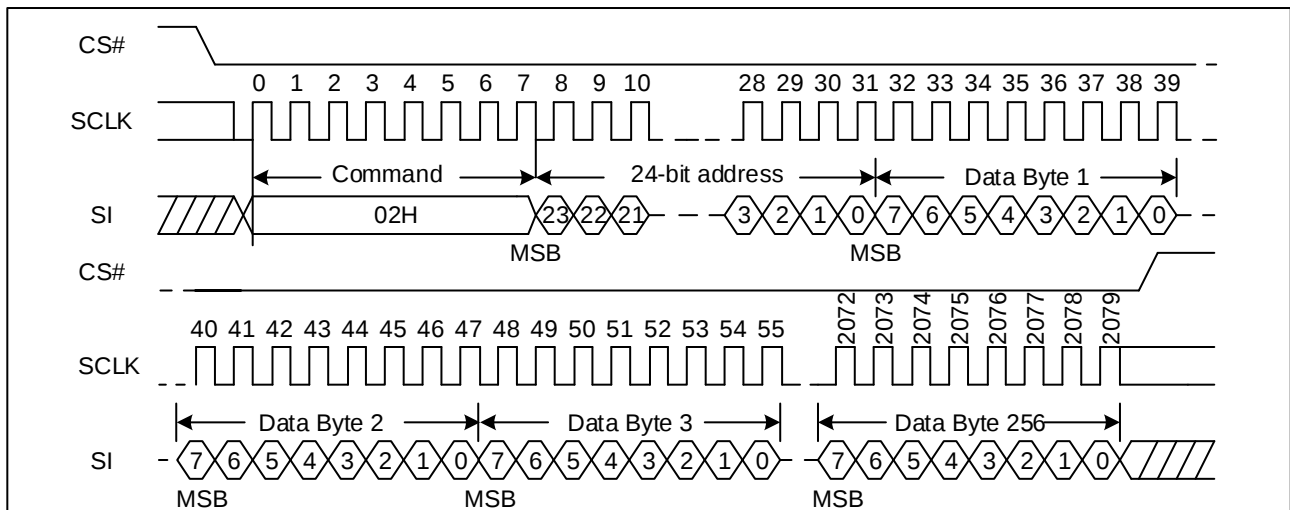


beyond the end of the current page are programmed from the start address of the same page (from the address whose 8 least significant bits (A7-A0) are all zero). CS# must be driven low for the entire duration of the sequence. The Page Program command sequence: CS# goes low → sending Page Program command → 3-byte address on SI → at least 1 byte data on SI → CS# goes high. If more than 256 bytes are sent to the device, previously latched data are discarded and the last 256 data bytes are guaranteed to be programmed correctly within the same page. If less than 256 data bytes are sent to device, they are correctly programmed at the requested addresses without having any effects on the other bytes of the same page. CS# must be driven high after the eighth bit of the last data byte has been latched in; otherwise the Page Program (PP) command is not executed.

As soon as CS# is driven high, the self-timed Page Program cycle (whose duration is t_{PP}) is initiated. While the Page Program cycle is in progress, the Status Register may be read to check the value of the Write in Progress (WIP) bit. The Write in Progress (WIP) bit is 1 during the self-timed Page Program cycle, and is 0 when it is completed. At some unspecified time before the cycle is completed, the Write Enable Latch (WEL) bit is reset.

A Page Program (PP) command applied to a page which is protected by the Block Protect (BP4, BP3, BP2, BP1, and BP0) is not executed.

Figure 13. Page Program Sequence Diagram



7.14 Quad Page Program (32H)

The Quad Page Program command is for programming the memory using four pins: IO0, IO1, IO2, and IO3. A Write Enable (WREN) command must previously have been executed to set the Write Enable Latch (WEL) bit before sending the Page Program command. The quad Page Program command is entered by driving CS# Low, followed by the command code (32H), three address bytes and at least one data byte on IO pins.

If more than 256 bytes are sent to the device, previously latched data are discarded and the last 256 data bytes are guaranteed to be programmed correctly within the same page. If less than 256 data bytes are sent to device, they are correctly programmed at the requested addresses without having any effects on the other bytes of the same page. CS# must be driven high after the eighth bit of the last data byte has been latched in; otherwise the Quad Page Program (PP) command is not executed.

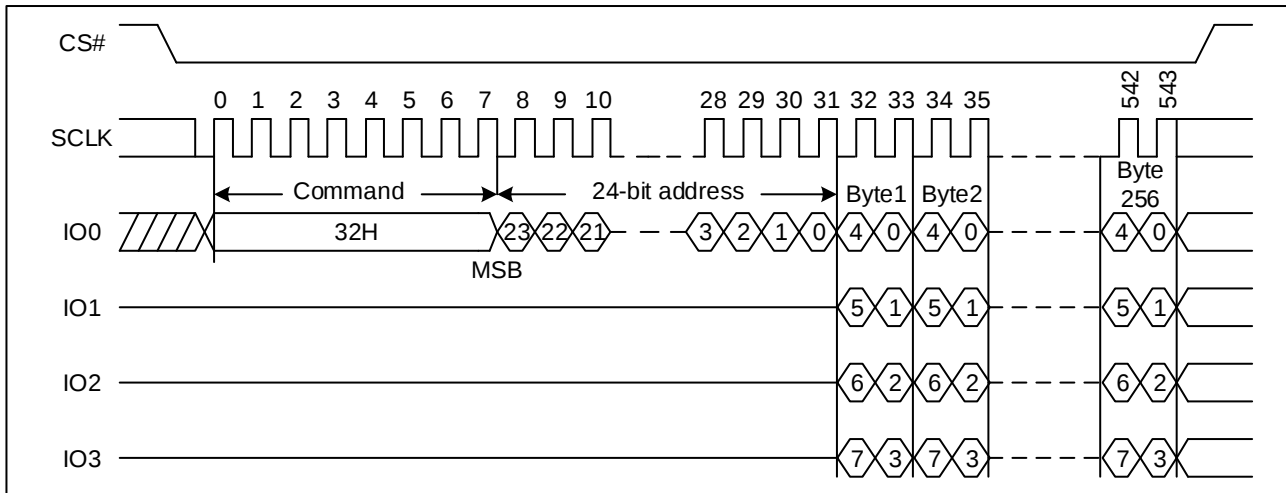
As soon as CS# is driven high, the self-timed Quad Page Program cycle (whose duration is t_{PP}) is initiated. While the Quad Page Program cycle is in progress, the Status Register may be read to check the value of the Write In Progress (WIP) bit. The Write in Progress (WIP) bit is 1 during the self-timed Quad Page Program cycle, and is 0 when it is completed. At some unspecified time before the cycle is completed, the Write Enable Latch (WEL) bit is reset.

A Quad Page Program command applied to a page which is protected by the Block Protect (BP4, BP3, BP2, BP1, and BP0)



is not executed.

Figure 14. Quad Page Program Sequence Diagram

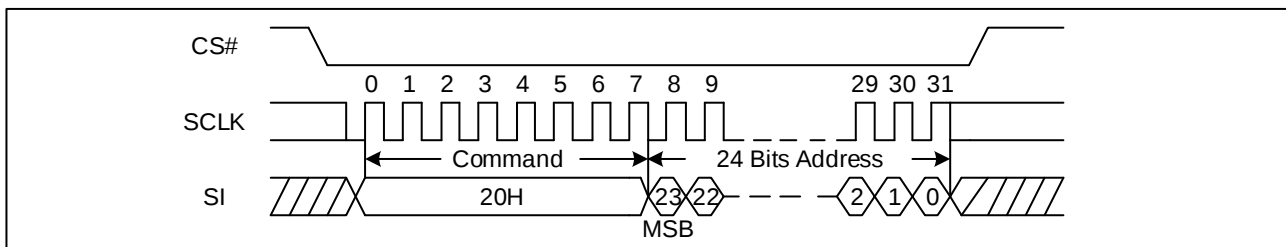


7.15 Sector Erase (SE) (20H)

The Sector Erase (SE) command is for erasing the all data of the chosen sector. A Write Enable (WREN) command must previously have been executed to set the Write Enable Latch (WEL) bit. The Sector Erase (SE) command is entered by driving CS# low, followed by the command code, and 3-address byte on SI. Any address inside the sector is a valid address for the Sector Erase (SE) command. CS# must be driven low for the entire duration of the sequence.

The Sector Erase command sequence: CS# goes low → sending Sector Erase command → 3-byte address on SI → CS# goes high. CS# must be driven high after the eighth bit of the last address byte has been latched in; otherwise the Sector Erase (SE) command is not executed. As soon as CS# is driven high, the self-timed Sector Erase cycle (whose duration is t_{SE}) is initiated. While the Sector Erase cycle is in progress, the Status Register may be read to check the value of the Write in Progress (WIP) bit. The Write in Progress (WIP) bit is 1 during the self-timed Sector Erase cycle, and is 0 when it is completed. At some unspecified time before the cycle is completed, the Write Enable Latch (WEL) bit is reset. A Sector Erase (SE) command applied to a sector which is protected by the Block Protect (BP4, BP3, BP2, BP1, and BP0) bit is not executed.

Figure 15. Sector Erase Sequence Diagram



7.16 32KB Block Erase (BE32) (52H)

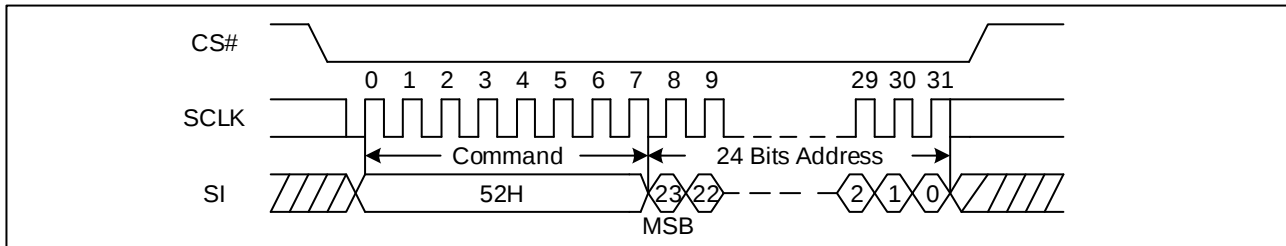
The 32KB Block Erase command is for erasing the all data of the chosen block. A Write Enable (WREN) command must previously have been executed to set the Write Enable Latch (WEL) bit. The 32KB Block Erase command is entered by driving CS# low, followed by the command code, and three address bytes on SI. Any address inside the block is a valid address for the 32KB Block Erase command. CS# must be driven low for the entire duration of the sequence.

The 32KB Block Erase command sequence: CS# goes low → sending 32KB Block Erase command → 3-byte address on



SI → CS# goes high. CS# must be driven high after the eighth bit of the last address byte has been latched in; otherwise the 32KB Block Erase command is not executed. As soon as CS# is driven high, the self-timed Block Erase cycle (whose duration is t_{BE1}) is initiated. While the Block Erase cycle is in progress, the Status Register may be read to check the value of the Write in Progress (WIP) bit. The Write in Progress (WIP) bit is 1 during the self-timed Block Erase cycle, and is 0 when it is completed. At some unspecified time before the cycle is completed, the Write Enable Latch (WEL) bit is reset. A 32KB Block Erase command applied to a block which is protected by the Block Protect (BP4, BP3, BP2, BP1, and BP0) bits is not executed.

Figure 16. 32KB Block Erase Sequence Diagram

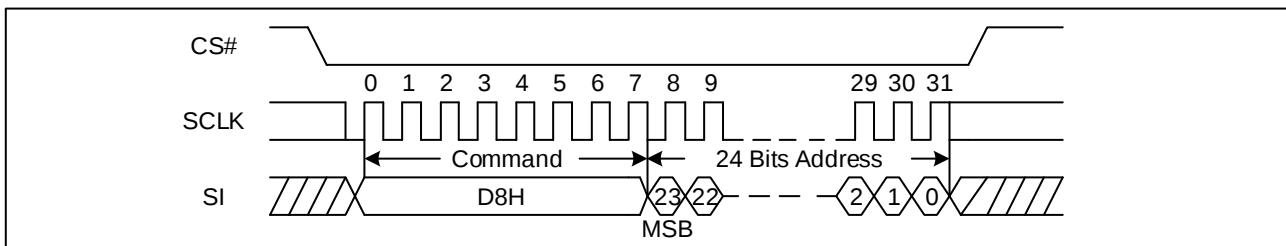


7.17 64KB Block Erase (BE64) (D8H)

The 64KB Block Erase command is for erasing the all data of the chosen block. A Write Enable (WREN) command must previously have been executed to set the Write Enable Latch (WEL) bit. The 64KB Block Erase command is entered by driving CS# low, followed by the command code, and three address bytes on SI. Any address inside the block is a valid address for the 64KB Block Erase command. CS# must be driven low for the entire duration of the sequence.

The 64KB Block Erase command sequence: CS# goes low → sending 64KB Block Erase command → 3-byte address on SI → CS# goes high. CS# must be driven high after the eighth bit of the last address byte has been latched in; otherwise the 64KB Block Erase command is not executed. As soon as CS# is driven high, the self-timed Block Erase cycle (whose duration is t_{BE2}) is initiated. While the Block Erase cycle is in progress, the Status Register may be read to check the value of the Write in Progress (WIP) bit. The Write in Progress (WIP) bit is 1 during the self-timed Block Erase cycle, and is 0 when it is completed. At some unspecified time before the cycle is completed, the Write Enable Latch (WEL) bit is reset. A 64KB Block Erase command applied to a block which is protected by the Block Protect (BP4, BP3, BP2, BP1, and BP0) bits is not executed.

Figure 17. 64KB Block Erase Sequence Diagram



7.18 Chip Erase (CE) (60/C7H)

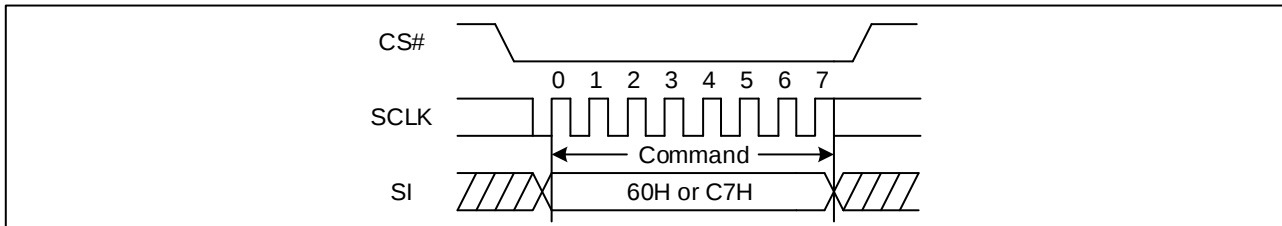
The Chip Erase (CE) command is for erasing the all data of the chip. A Write Enable (WREN) command must previously have been executed to set the Write Enable Latch (WEL) bit. The Chip Erase (CE) command is entered by driving CS# Low, followed by the command code on Serial Data Input (SI). CS# must be driven Low for the entire duration of the sequence.

The Chip Erase command sequence: CS# goes low → sending Chip Erase command → CS# goes high. CS# must be



driven high after the eighth bit of the command code has been latched in; otherwise the Chip Erase command is not executed. As soon as CS# is driven high, the self-timed Chip Erase cycle (whose duration is t_{CE}) is initiated. While the Chip Erase cycle is in progress, the Status Register may be read to check the value of the Write in Progress (WIP) bit. The Write in Progress (WIP) bit is 1 during the self-timed Chip Erase cycle, and is 0 when it is completed. At some unspecified time before the cycle is completed, the Write Enable Latch (WEL) bit is reset. The Chip Erase (CE) command is executed, if the Block Protect (BP2, BP1, and BP0) bits are 0 and CMP=0 or the Block Protect (BP2, BP1, and BP0) bits are 1 and CMP=1. The Chip Erase (CE) command is ignored if one or more sectors are protected.

Figure 18. Chip Erase Sequence Diagram

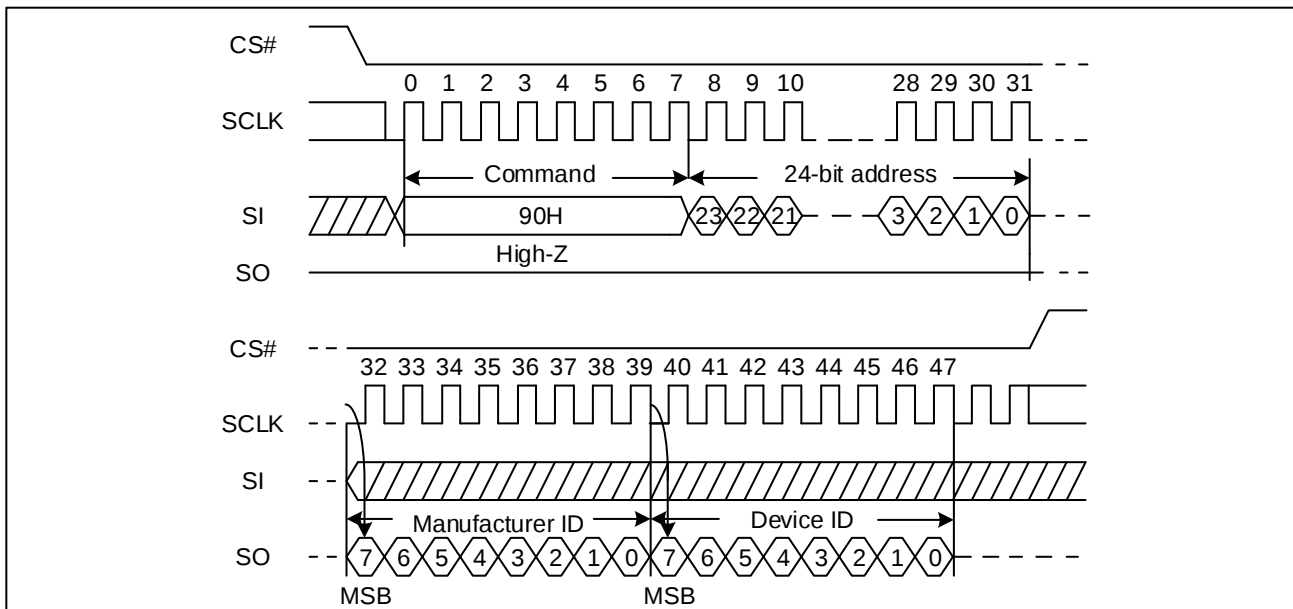


7.19 Read Manufacture ID/ Device ID (REMS) (90H)

The Read Manufacturer/Device ID command is an alternative to the Release from Power-Down / Device ID command that provides both the JEDEC assigned Manufacturer ID and the specific Device ID.

The command is initiated by driving the CS# pin low and shifting the command code "90H" followed by a 24-bit address (A23-A0) of 000000H. After which, the Manufacturer ID and the Device ID are shifted out on the falling edge of SCLK with most significant bit (MSB) first.

Figure 19. Read Manufacture ID/ Device ID Sequence Diagram



7.20 Read Identification (RDID) (9FH)

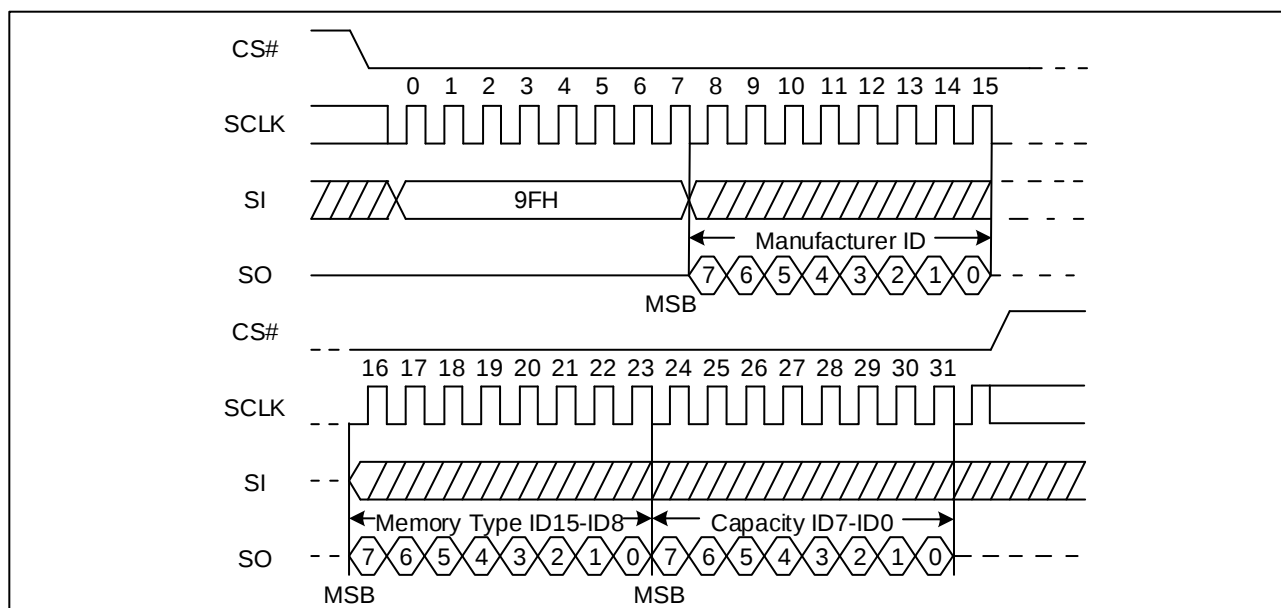
The Read Identification (RDID) command allows the 8-bit manufacturer identification to be read, followed by two bytes of device identification. The device identification indicates the memory type in the first byte, and the memory capacity of the device in the second byte. The Read Identification (RDID) command while an Erase or Program cycle is in progress, is not decoded, and has no effect on the cycle that is in progress. The Read Identification (RDID) command should not be issued



while the device is in Deep Power-Down Mode.

The device is first selected by driving CS# low. Then, the 8-bit command code for the command is shifted in. This is followed by the 24-bit device identification, stored in the memory. Each bit is shifted out on the falling edge of Serial Clock. The Read Identification (RDID) command is terminated by driving CS# high at any time during data output. When CS# is driven high, the device is in the Standby Mode. Once in the Standby Mode, the device waits to be selected, so that it can receive, decode and execute commands.

Figure 20. Read Identification ID Sequence Diagram

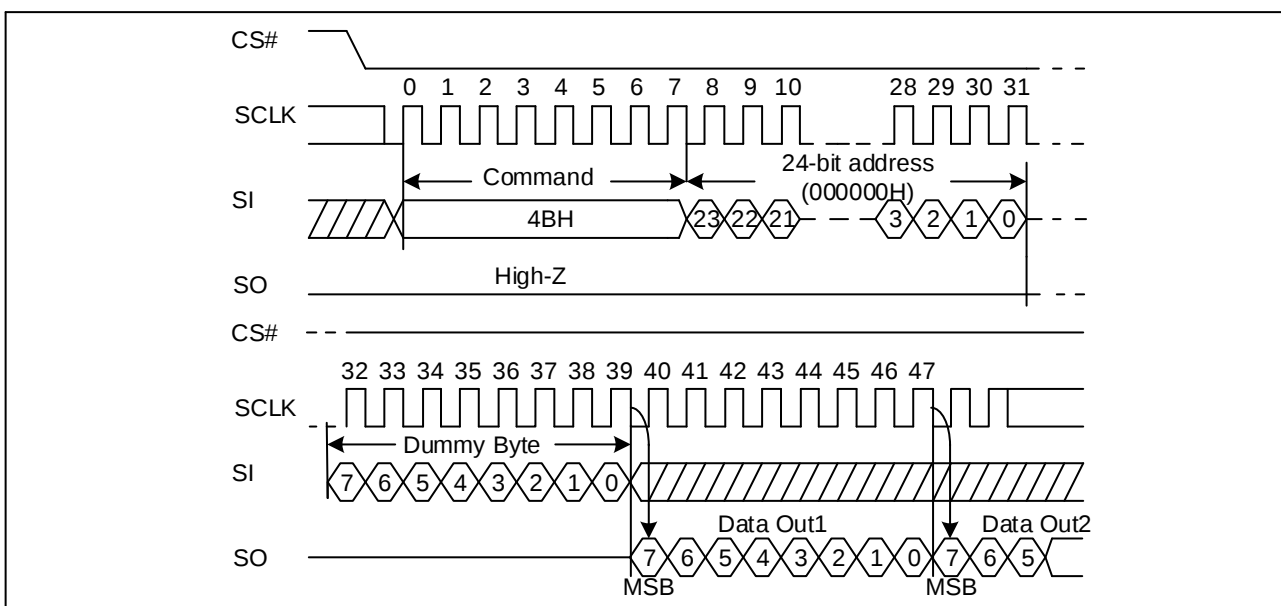


7.21 Read Unique ID (4BH)

The Read Unique ID command accesses a factory-set read-only 128bit number that is unique to each device. The Unique ID can be used in conjunction with user software methods to help prevent copying or cloning of a system.

The Read Unique ID command sequence: CS# goes low → sending Read Unique ID command → 3-Byte Address (000000H) → Dummy Byte → 128bit Unique ID Out → CS# goes high.

Figure 21. Read Unique ID Sequence Diagram





7.22 Erase Security Registers (44H)

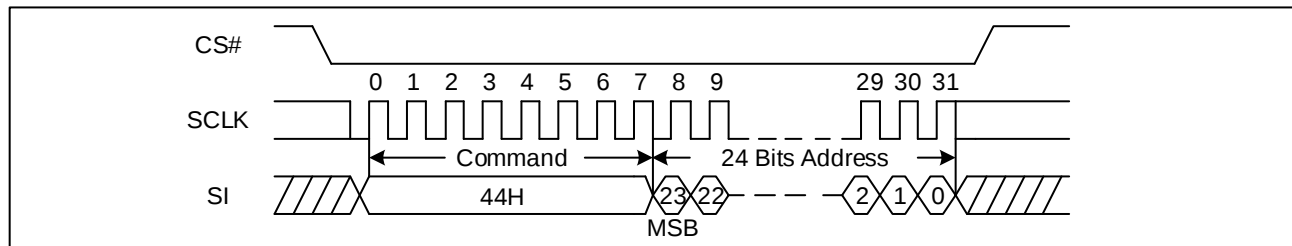
The GD25R128E provides 3x1024-Byte Security Registers which can be erased and programmed individually. These registers may be used by the system manufacturers to store security and other important information separately from the main memory array.

The Erase Security Registers command is similar to Sector/Block Erase command. A Write Enable (WREN) command must previously have been executed to set the Write Enable Latch (WEL) bit.

The Erase Security Registers command sequence: CS# goes low → sending Erase Security Registers command → 3-byte address on SI → CS# goes high. The command sequence is shown below. CS# must be driven high after the eighth bit of the last address byte has been latched in; otherwise the Erase Security Registers command is not executed. As soon as CS# is driven high, the self-timed Erase Security Registers cycle (whose duration is t_{SE}) is initiated. While the Erase Security Registers cycle is in progress, the Status Register may be read to check the value of the Write in Progress (WIP) bit. The Write in Progress (WIP) bit is 1 during the self-timed Erase Security Registers cycle, and is 0 when it is completed. At some unspecified time before the cycle is completed, the Write Enable Latch (WEL) bit is reset. The Security Registers Lock Bit (LB1, LB2, LB3) in the Status Register can be used to OTP protect the security registers. Once the LB bit is set to 1, the Security Registers will be permanently locked; the Erase Security Registers command will be ignored.

Address	A23-16	A15-12	A11-10	A9-0
Security Register #1	00H	0001b	00b	Don't care
Security Register #2	00H	0010b	00b	Don't care
Security Register #3	00H	0011b	00b	Don't care

Figure 22. Erase Security Registers command Sequence Diagram



7.23 Program Security Registers (42H)

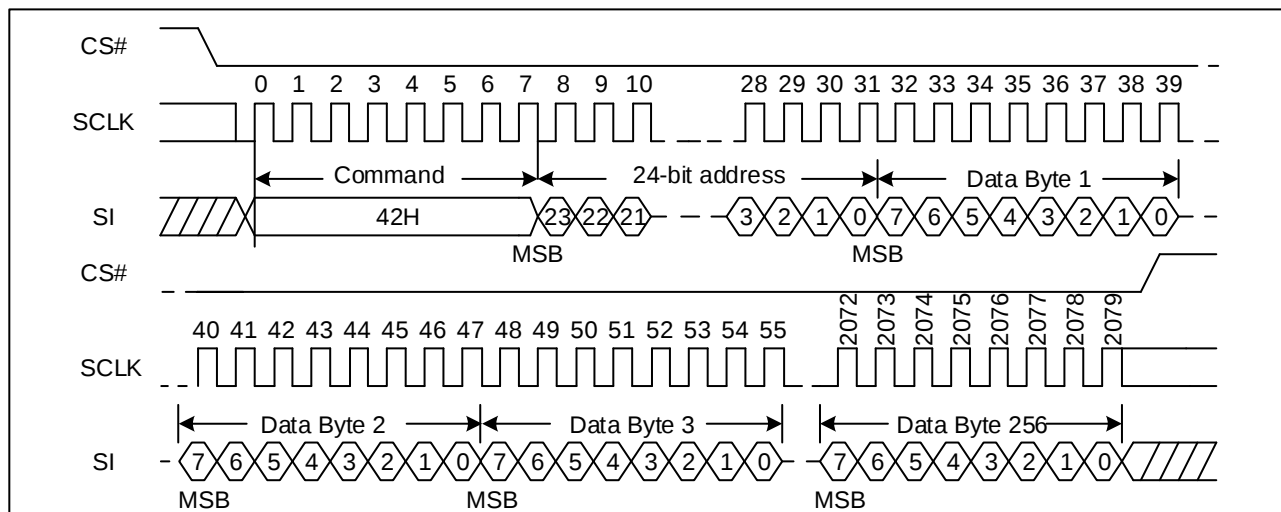
The Program Security Registers command is similar to the Page Program command. Each security register contains four pages content. A Write Enable (WREN) command must previously have been executed to set the Write Enable Latch (WEL) bit before sending the Program Security Registers command. The Program Security Registers command is entered by driving CS# Low, followed by the command code (42H), three address bytes and at least one data byte on SI. As soon as CS# is driven high, the self-timed Program Security Registers cycle (whose duration is t_{PP}) is initiated. While the Program Security Registers cycle is in progress, the Status Register may be read to check the value of the Write in Progress (WIP) bit. The Write in Progress (WIP) bit is 1 during the self-timed Program Security Registers cycle, and is 0 when it is completed. At some unspecified time before the cycle is completed, the Write Enable Latch (WEL) bit is reset.

If the Security Registers Lock Bit (LB1, LB2, LB3) is set to 1, the Security Registers will be permanently locked. Program Security Registers command will be ignored.

Address	A23-16	A15-12	A11-10	A9-0
Security Register #1	00H	0001b	00b	Byte Address
Security Register #2	00H	0010b	00b	Byte Address
Security Register #3	00H	0011b	00b	Byte Address



Figure 23. Program Security Registers command Sequence Diagram

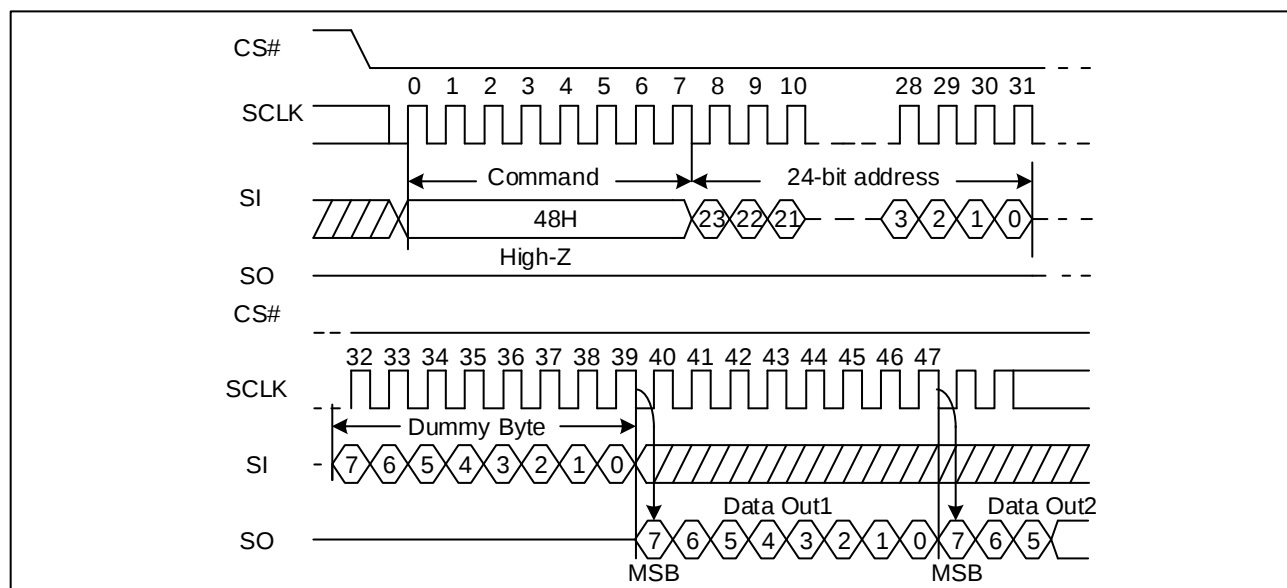


7.24 Read Security Registers (48H)

The Read Security Registers command is similar to Fast Read command. The command is followed by a 3-byte address (A23-A0) and a dummy byte, and each bit is latched-in on the rising edge of SCLK. Then the memory content, at that address, is shifted out on SO, and each bit is shifted out, at a Max frequency f_c , on the falling edge of SCLK. The first byte addressed can be at any location. The address is automatically incremented to the next higher address after each byte of data is shifted out. Once the A9-0 address reaches the last byte of the register (Byte 3FFH), it will reset to 000H, the command is completed by driving CS# high.

Address	A23-16	A15-12	A11-10	A9-0
Security Register #1	00H	0001b	00b	Byte Address
Security Register #2	00H	0010b	00b	Byte Address
Security Register #3	00H	0011b	00b	Byte Address

Figure 24. Read Security Registers command Sequence Diagram



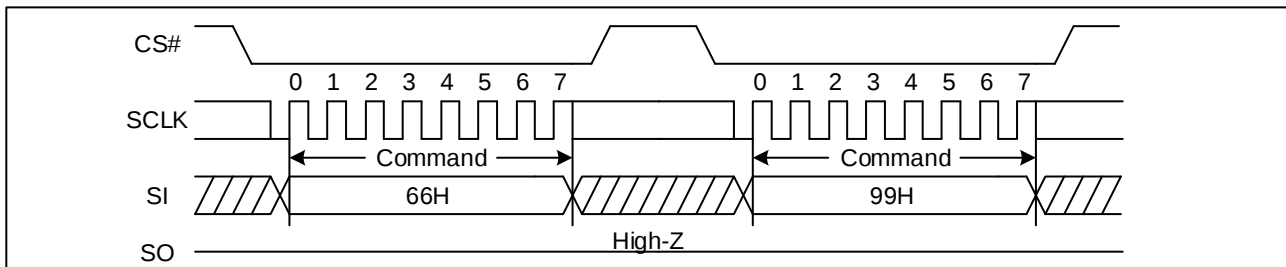


7.25 Enable Reset (66H) and Reset (99H)

If the Reset command is accepted, any on-going internal operation will be terminated and the device will return to its default power-on state and lose all the current volatile settings, such as Volatile Status Register bits, Write Enable Latch status (WEL), Program/Erase Suspend status, Read Parameter setting (P7-P0), Deep Power Down Mode, Continuous Read Mode bit setting (M7-M0) and Wrap Bit Setting (W6-W4).

The “Enable Reset (66H)” and “Reset (99H)” command sequence as follow: CS# goes low → Sending Enable Reset command → CS# goes high → CS# goes low → Sending Reset command → CS# goes high. Once the Reset command is accepted by the device, the device will take approximately t_{RST} / t_{RST_E} to reset. During this period, no command will be accepted. Data corruption may happen if there is an on-going or suspended internal Erase or Program operation when Reset command sequence is accepted by the device. It is recommended to check the WIP bit and the SUS1/SUS2 bit in Status Register before issuing the Reset command sequence.

Figure 25. Enable Reset and Reset command Sequence Diagram

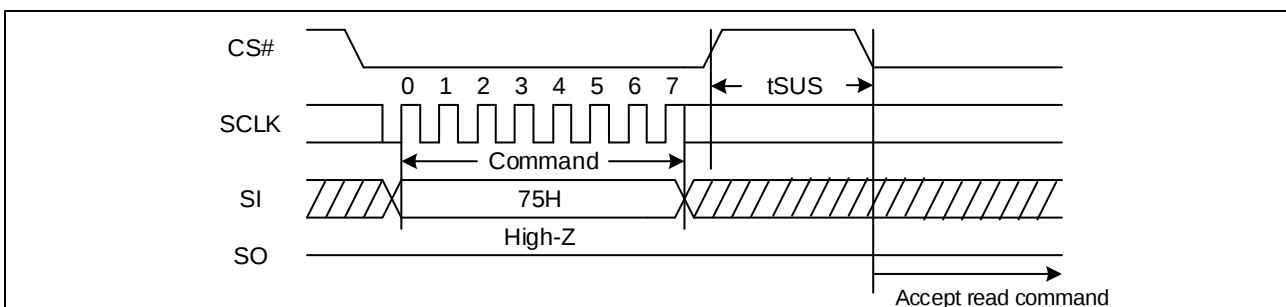


7.26 Program/Erase Suspend (PES) (75H)

The Program/Erase Suspend command “75H”, allows the system to interrupt a page program or sector/block erase operation and then read data from any other sector or block. The Write Status Register command (01H, 31H, 11H) and Erase/Program Security Registers command (44H, 42H) and Erase commands (20H, 52H, D8H, C7H, 60H) and Page Program command (02H, 32H) are not allowed during Program suspend. The Write Status Register command (01H, 31H, 11H) and Erase Security Registers command (44H) and Erase commands (20H, 52H, D8H, C7H, 60H) are not allowed during Erase suspend. Program/Erase Suspend is valid only during the page program or sector/block erase operation. A maximum of time of “ t_{sus} ” (See AC Characteristics) is required to suspend the program/erase operation.

The Program/Erase Suspend command will be accepted by the device only if the SUS1/SUS2 bit in the Status Register equal to 0 and WIP bit equal to 1 while a Page Program or a Sector or Block Erase operation is on-going. If the SUS1/SUS2 bit equal to 1 or WIP bit equal to 0, the Suspend command will be ignored by the device. The WIP bit will be cleared from 1 to 0 within “ t_{sus} ” and the SUS1/SUS2 bit will be set from 0 to 1 immediately after Program/Erase Suspend. A power-off during the suspend period will reset the device and release the suspend state.

Figure 26. Program/Erase Suspend Sequence Diagram

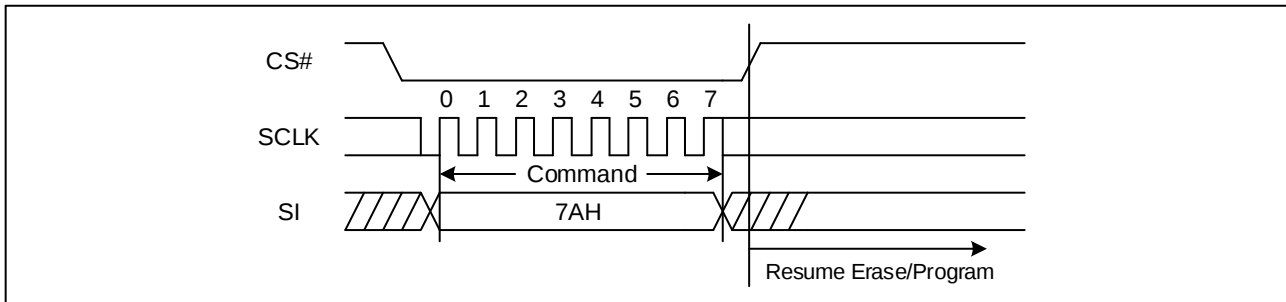




7.27 Program/Erase Resume (PER) (7AH)

The Program/Erase Resume command must be written to resume the program or sector/block erase operation after a Program/Erase Suspend command. The Program/Erase Resume command will be accepted by the device only if the SUS1/SUS2 bit equal to 1 and the WIP bit equal to 0. After issued the SUS1/SUS2 bit in the status register will be cleared from 1 to 0 immediately, the WIP bit will be set from 0 to 1 within 200ns and the Sector or Block will complete the erase operation or the page will complete the program operation. The Program/Erase Resume command will be ignored unless a Program/Erase Suspend is active.

Figure 27. Program/Erase Resume Sequence Diagram



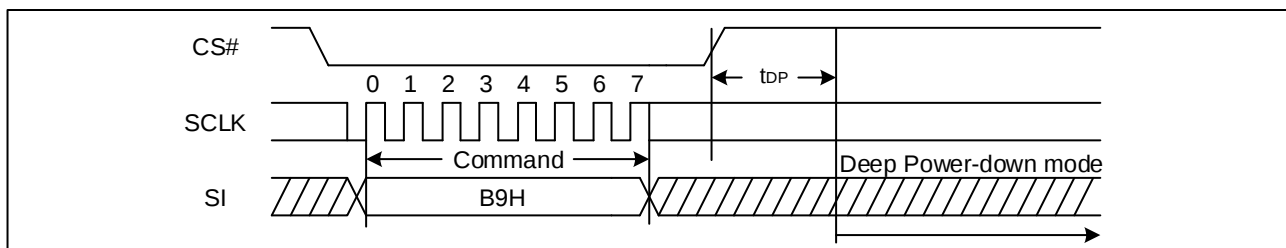
7.28 Deep Power-Down (DP) (B9H)

Executing the Deep Power-Down (DP) command is the only way to put the device in the lowest consumption mode (the Deep Power-Down Mode). It can also be used as an extra software protection mechanism, while the device is not in active use, since in this mode, the device ignores all Write, Program and Erase commands. Driving CS# high deselects the device, and puts the device in the Standby Mode (if there is no internal cycle currently in progress). But this mode is not the Deep Power-Down Mode. The Deep Power-Down Mode can only be entered by executing the Deep Power-Down (DP) command. Once the device has entered the Deep Power-Down Mode, all commands are ignored except the Release from Deep Power-Down and Read Device ID (RDI) command or software reset command. The Release from Deep Power-Down and Read Device ID (RDI) command releases the device from Deep Power-Down mode, also allows the Device ID of the device to be output on SO.

The Deep Power-Down Mode automatically stops at Power-Down, and the device always in the Standby Mode after Power-Up.

The Deep Power-Down command sequence: CS# goes low → sending Deep Power-Down command → CS# goes high. CS# must be driven high after the eighth bit of the command code has been latched in; otherwise the Deep Power-Down (DP) command is not executed. As soon as CS# is driven high, it requires a delay of t_{DP} before the supply current is reduced to I_{CC2} and the Deep Power-Down Mode is entered. Any Deep Power-Down (DP) command, while an Erase, Program or Write cycle is in progress, is rejected without having any effects on the cycle that is in progress.

Figure 28. Deep Power-Down Sequence Diagram





7.29 Release from Deep Power-Down and Read Device ID (RDI) (ABH)

The Release from Power-Down and Read Device ID command is a multi-purpose command. It can be used to release the device from the Power-Down state or obtain the device's electronic identification (ID) number.

To release the device from the Power-Down state, the command is issued by driving the CS# pin low, shifting the instruction code "ABH" and driving CS# high. Release from Power-Down will take the time duration of t_{RES1} (See AC Characteristics) before the device will resume normal operation and other commands are accepted. The CS# pin must remain high during the t_{RES1} time duration.

When used only to obtain the Device ID while not in the Power-Down state, the command is initiated by driving the CS# pin low and shifting the instruction code "ABH" followed by 3-dummy byte. The ID7-ID0 are then shifted out on the falling edge of SCLK with most significant bit (MSB) first. The Device ID value is listed in Manufacturer and Device Identification table. The Device ID can be read continuously. The command is completed by driving CS# high.

When used to release the device from the Power-Down state and obtain the ID7-ID0, the command is the same as previously described, except that after CS# is driven high it must remain high for a time duration of t_{RES2} (See AC Characteristics). After this time duration the device will resume normal operation and other commands will be accepted. If the Release from Power-Down / Device ID command is issued while an Erase, Program or Write cycle is in process (when WIP equals 1) the command is ignored and will not have any effects on the current cycle.

Figure 29. Release Power-Down Sequence Diagram

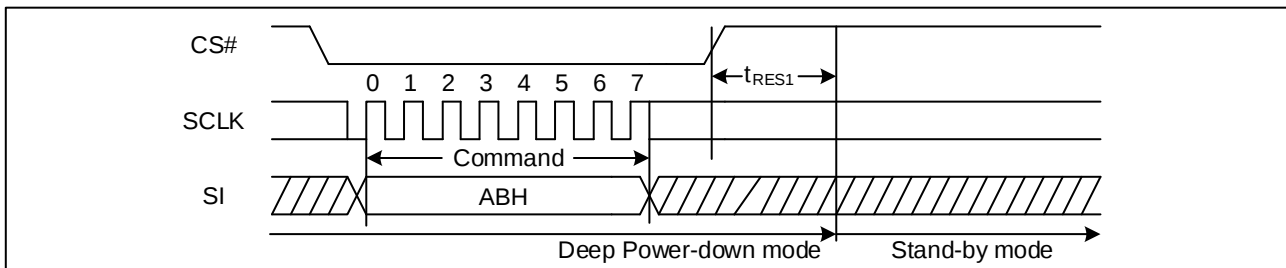
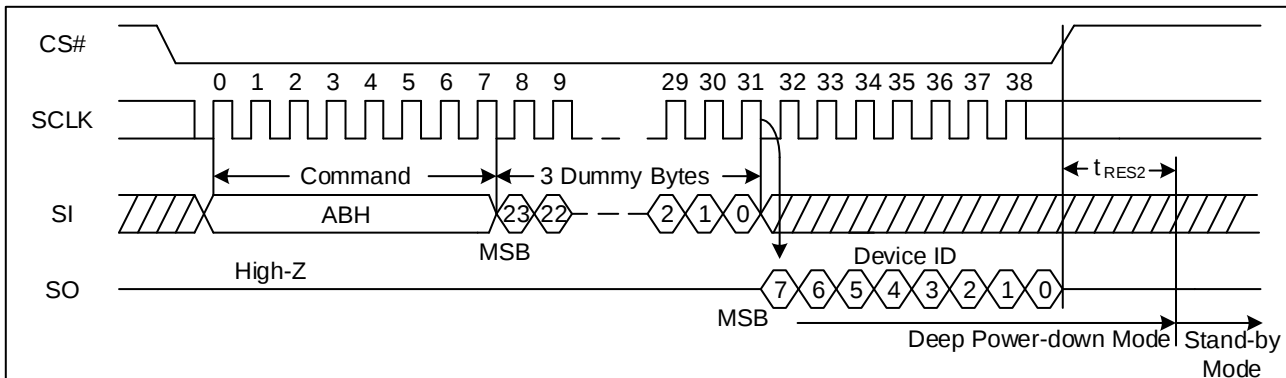


Figure 30. Release Power-Down/Read Device ID Sequence Diagram



7.30 Read Serial Flash Discoverable Parameter (5AH)

The Serial Flash Discoverable Parameter (SFDP) standard provides a consistent method of describing the functional and feature capabilities of serial flash devices in a standard set of internal parameter tables. These parameter tables can be interrogated by host system software to enable adjustments needed to accommodate divergent features from multiple vendors. The concept is similar to the one found in the Introduction of JEDEC Standard, JESD68 on CFI. SFDP is a standard of JEDEC Standard No.216B.



Figure 31. Read Serial Flash Discoverable Parameter command Sequence Diagram

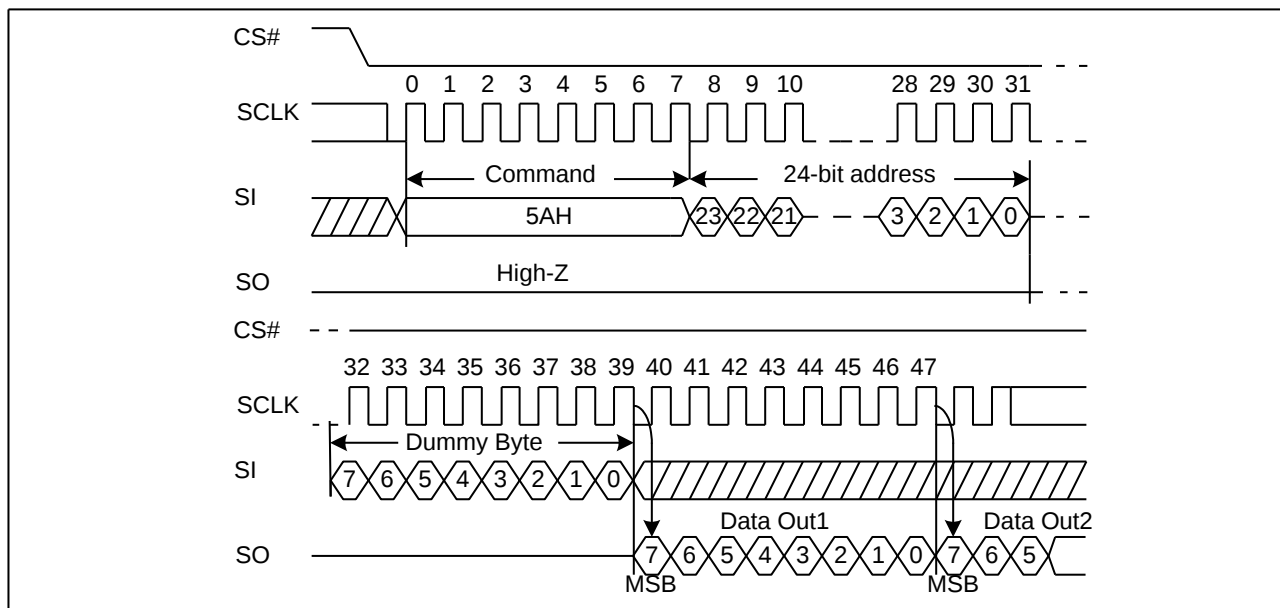


Table 9. Signature and Parameter Identification Data Values (Please contact GigaDevice for Details)



8 RPMC COMMANDS DESCRIPTION

Table 10. Replay Protected Monotonic Counter (RPMC) Commands

Function	Opcode Phase 8 bits	Payload Phase Max 512 Bits		Comment
		Byte#	Field Description	
Command: Write Root Key Register	OP1	1 2 3 4-35 36-63	CmdType[7:0] = 00H CounterAddr[7:0] = 8 Bits Reserved[7:0] = 8 Bits RootKey[255:0] = 256 Bits TruncatedSign[223:0] = 224 Bits	OP1 + Payload phase driven by host controller. Root Key Register is written only once.
Command: Update HMAC Key Register	OP1	1 2 3 4-7 8-39	CmdType[7:0] = 01H CounterAddr[7:0] = 8 Bits Reserved[7:0] = 8 Bits KeyData[31:0] = 32 Bits, Signature[255:0] = 256 Bits	OP1 + Payload phase is Issued by host controller on every power up to initialize HMAC Key Register.
Command: Increment Monotonic Counter	OP1	1 2 3 4-7 8-39	CmdType[7:0] = 02H CounterAddr[7:0] = 8 Bits Reserved[7:0] = 8 Bits CounterData[31:0] = 32 Bits, Signature[255:0] = 256 Bits	OP1 + Payload Phase is Issued by host controller during runtime to increment the counter.
Command: Request Monotonic Counter	OP1	1 2 3 4-15 16-47	CmdType[7:0] = 03H CounterAddr[7:0] = 8 Bits Reserved[7:0] = 8 Bits Tag [95:0] = 96 Bits Signature[255:0] = 256 Bits	OP1 + Payload Phase is Issued by host controller during runtime to request counter data
Command: Read Data	OP2	2 3-14 15-18 19-50	ExtendedStatus[7:0] = 8 Bits Tag[95:0] = 96 Bits CounterData[31:0] = 32 Bits Signature[255:0] = 256 Bits	OP2 is issued by Host Controller generally after an OP1. SPI Flash device responds with the Payload phase to return Extended Status and counter data.

All individual fields are Byte wide fields. For a multi-byte field, Most Significant Byte is issued first; Least Significant Byte is issued last. Within a Byte, Most Significant Bit is issued first; Least Significant Bit is issued last. CmdType is always the first byte issued after OP1 commands. OP2 delay is the same as Fast Read Command delay which is 8 dummy bits.

Table 11. OP1 and OP2 are defined for 1-1-1 mode.

Byte #	0	1	2	3	4	5	6	...	..	..
Name	OP1	CmdType	Counter Address	As defined in the table above						
Name	OP2	8 Dummy clocks	Extended Status[7:0]	As defined in the table above						

After an OP1 command is received, the SPI Flash will indicate status busy indication using either the status register or extended status register as defined below.



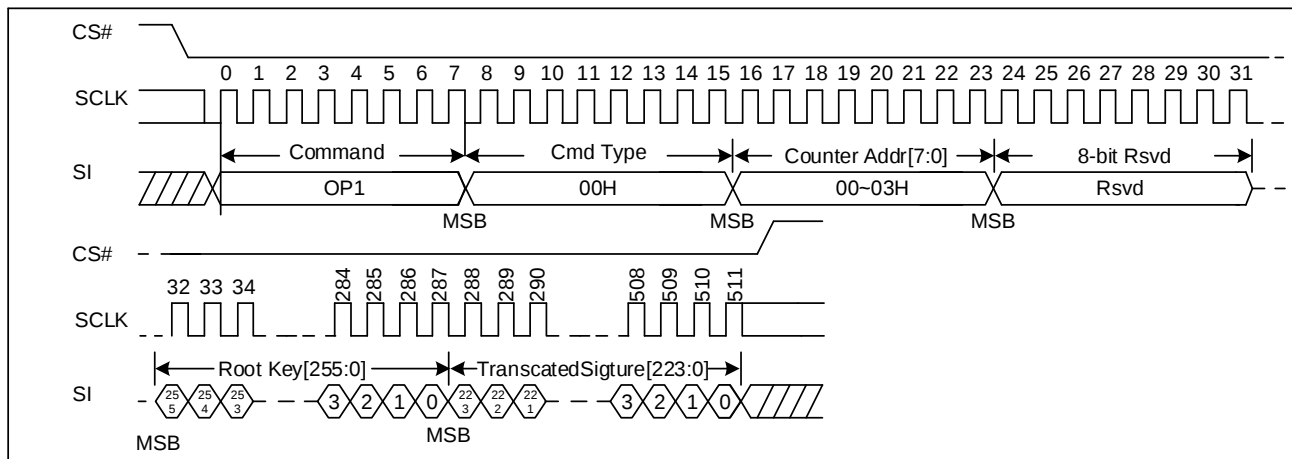
Table 12. Extended Status Register Definition

Extended Status[7:0]	Applicable CmdType(s)	Description
00000000	-	Power On State (OP2 issued directly after power-up).
10000000	00,01,02,03	This status is set on successful completion (no errors) of OP1 command.
0xxxxxx1	00,01,02,03,04-0FF	This bit is set to 1, when device is busy executing OP1 command. It is reset to 0 when OP1 command execution is done.
0xxxxx1x	00,01	This bit is set only when the correct payload size is received. When cmdtype = 0, this error bit must be set on Root Key Register Overwrite or Counter Address out of range or Truncated Signature mis-match error. For cmdtype = 01 this bit is set when the corresponding monotonic counter is uninitialized
0xxxx1xx	00,01,02,03	This bit is set on Signature Mismatch, Counter Address out of range when correct payload size is received; or Cmdtype is out of range; or incorrect payload size is received.
0xxx1xxx	02,03	This bit is set on HMAC Key Register or monotonic counter uninitialized on previous OP1 command when correct payload size is received
0xx1xxxx	02	This bit is set on Counter Data Mismatch on previous increment when correct payload size is received
0x1xxxxx	-	Fatal Error. It is set when no valid counter is found after initialization.
Current value		Extended status register will not be updated until first 8 bits of OP1 is received. The correct error type is reflected for any OP1 operation that exceeds a minimum of 16 clocks with active chip-select.

8.1 Command: Write Root Key Register

This command is used to initialize the Root Key Register corresponding to the received Counter Address with the received Root Key. It is suggested to be used in an OEM manufacturing environment when the SPI Flash Controller and SPI Flash are powered together for the first time.

Figure 32. Write Root Key Register Sequence Diagram



Truncated signature field is the same as least significant 224 bits of HMAC-SHA-256 based signature computed based on



received input parameters:

- HMAC message[31:0] = (OpCode[7:0], CmdType[7:0], CounterAddr[7:0], Reserved[7:0])
- HMAC Key[255:0] = Root_Key[255:0]

If Root Key != 256'HFF..FF then this command can be executed one time. If the received transaction is error free SPI Flash device successfully executes the command and posts "successful completion" extended status.

Root Key Register Write with root key is = 256'HFF...FF can be used as a temporary key.

If the received transaction has errors the SPI Flash does not execute the transaction and posts the corresponding error in extended status.

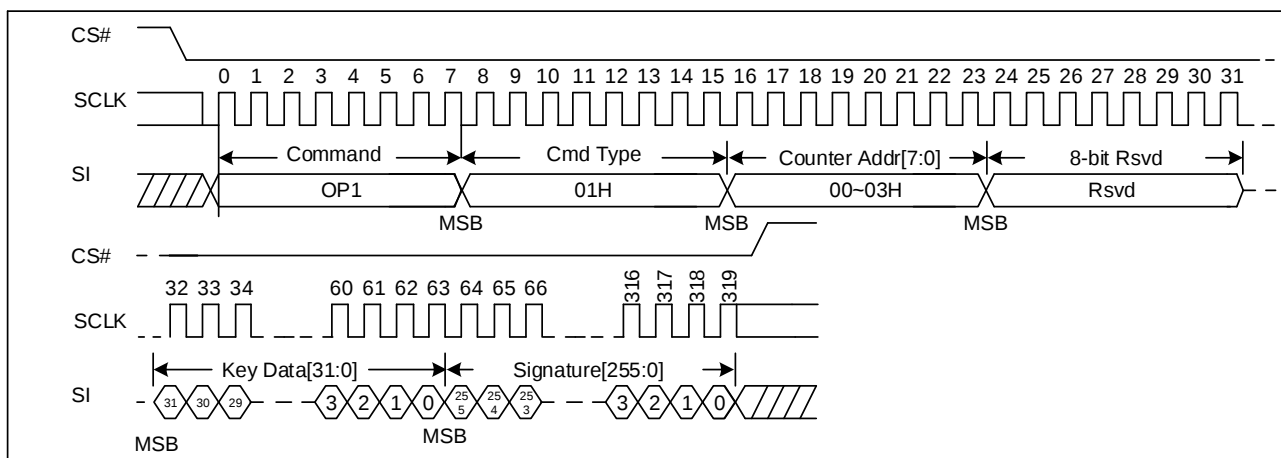
Table 13. Expected Extended Status [7:0] results

Extended Status[7:0]	Applicable CmdType(s)	Description
10000000	00	Successful completion
0xxxxxx1	00	If Busy_Polling_Method bit in SFDP table is zero, then this bit must be set to 1, when device is busy executing command. It is reset to 0 when OP1 command execution is done. If Busy_Polling_Method bit in SFDP table is one, then this bit is ignored by the controller.
0xxxxx1x	00	This bit is only set when correct payload size is received. It is set on Root Key Register Overwrite or Counter Address is out of range or when there is a truncated signature mismatch error
0xxxx1xx	00	This bit is set when incorrect payload size is received.

8.2 Command: Update HMAC Key Register

This command is used by the SPI Flash Controller to update the HMAC-Key register corresponding to the received Counter Address with a new HMAC key calculated based on received input. This command must be issued on every power cycle event on the interface. The HMAC key storage is volatile.

Figure 33. Update HMAC Key Register Sequence Diagram



Signature matches the HMAC-SHA-256 based signature computed based on received input parameters. This command performs two HMAC-SHA-256 operations.

- HMAC-SHA-256 Operation 1 Output = HMAC_Storage[255:0]
 - HMAC Message[31:0] = KeyData[31:0]
 - HMAC Key[255:0] = Root_Key_Register[CounterAddr][255:0]



- HMAC-SHA-256 Operation 2 Output = HMAC-SHA-256 based signature[255:0]
 - HMAC message[63:0] = (OpCode[7:0], CmdType[7:0], CounterAddr[7:0], Reserved[7:0], KeyData[31:0])
 - HMAC Key[255:0] = HMAC_Storage[255:0]

If the received transaction is error free SPI Flash device successfully executes the command and posts “successful completion” extended status.

If the received transaction has errors the SPI Flash does not execute the transaction and posts the corresponding error in extended status.

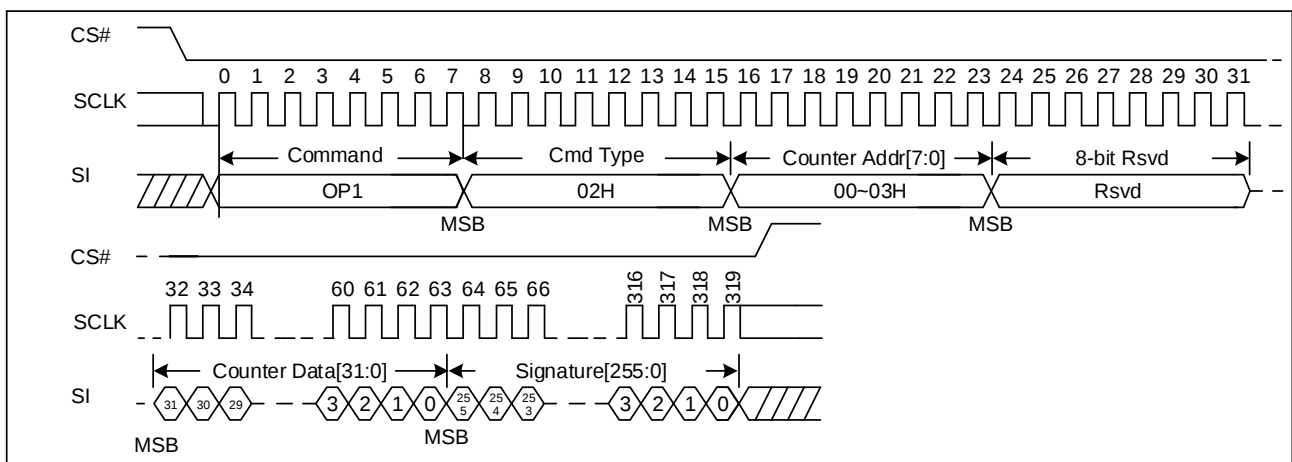
Table 14. Expected Extended Status [7:0] results

Extended Status[7:0]	Applicable CmdType(s)	Description
10000000	01	This status is set on successful completion (no errors) of OP1 command.
0xxxxx1	01	This bit is set to 1, when device is busy executing OP1 command. It is reset to 0 when OP1 command execution is done.
0xxxx1x	01	This bit is set only when the correct payload size is received. This bit is set when the corresponding monotonic counter is uninitialized
0xxxx1xx	01	This bit is set on Signature Mismatch, Counter Address out of range when correct payload size is received; or incorrect payload size is received.

8.3 Command: Increment Monotonic Counter

This command is used by the SPI Flash Controller to increment the Monotonic counter by 1 inside the SPI Flash Device.

Figure 34. Increment Monotonic Counter Sequence Diagram



The requested Signature matches the HMAC-SHA-256 based signature computed based on received input parameters.

The received Counter Data matches the current value of the counter read from the SPI Flash.

- HMAC Message[63:0] = (OpCode[7:0], CmdType[7:0], CounterAddr[7:0], Reserved[7:0], CounterData[31:0])
- HMAC Key[255:0] = HMAC_Key_Register [Counter_Address][255:0]

If the received transaction is error free SPI Flash device successfully executes the command and posts “successful completion” extended status. If the received transaction has errors the SPI Flash does not execute the transaction and posts the corresponding error in extended status.



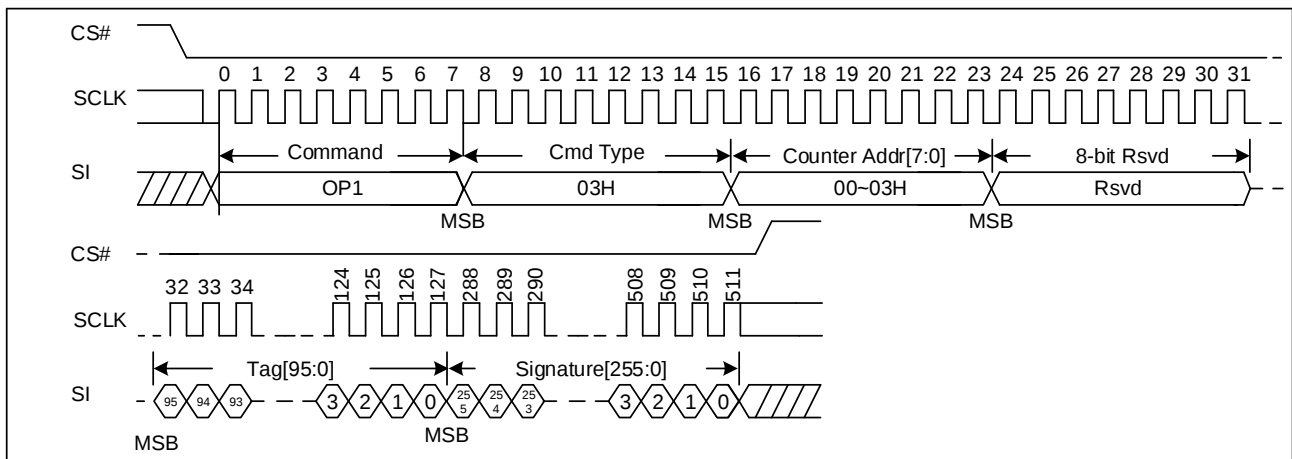
Table 15. Expected Extended Status [7:0] results

Extended Status[7:0]	Applicable CmdType(s)	Description
10000000	02	This status is set on successful completion (no errors) of OP1 command.
0xxxxxx1	02	This bit is set to 1, when device is busy executing OP1 command. It is reset to 0 when OP1 command execution is done.
0xxxx1xx	02	This bit is set on Signature Mismatch, Counter Address out of range when correct payload size is received; or incorrect payload size is received.
0xxx1xxx	02	This bit is set only when the correct payload size is received. This bit must be set on HMAC Key Register or Monotonic Counter is uninitialized on previous OP1 command.
0xx1xxxx	02	This bit is set only when the correct payload size is received. The bit must be set when the received counter data filed does not match the actual counter value read from the SPI Flash device.

8.4 Command: Request Monotonic Counter

This command is used by the SPI Flash Controller to request the Monotonic counter value inside the SPI Flash Device.

Figure 35. Request Monotonic Counter Sequence Diagram



The requested Signature matches the HMAC-SHA-256 based signature computed based on received input parameters.

- HMAC Message[127:0] = (OpCode[7:0], CmdType[7:0], CounterAddr[7:0], Reserved[7:0], Tag[95:0])
- HMAC Key[255:0] = HMAC_Key_Register[Counter_Address][255:0]

If the received transaction is error free SPI Flash device successfully executes the command and posts “successful completion” extended status. If the received transaction has errors the SPI Flash does not execute the transaction and posts the corresponding error in extended status.



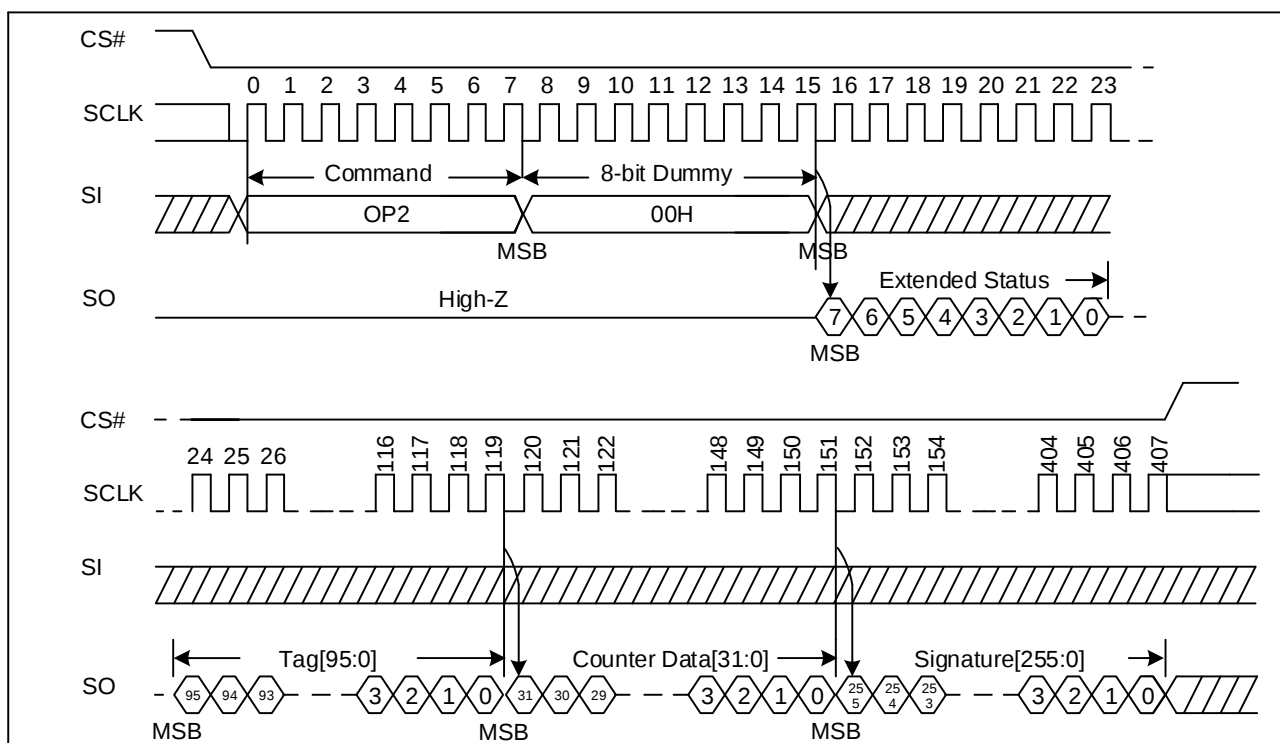
Table 16. Expected Extended Status [7:0] results

Extended Status[7:0]	Applicable CmdType(s)	Description
10000000	03	This status is set on successful completion (no errors) of OP1 command.
0xxxxxx1	03	This bit is set to 1, when device is busy executing OP1 command. It is reset to 0 when OP1 command execution is done.
0xxxx1xx	03	This bit is set on Signature Mismatch, Counter Address out of range when correct payload size is received; or Cmdtype is out of range; or incorrect payload size is received.
0xxx1xxx	03	This bit is set only when the correct payload size is received. This bit must be set on HMAC Key Register or Monotonic Counter is uninitialized on previous OP1 command.

8.5 Command: Read Data

This command is used by the SPI Flash Controller to read extended status from any previously issued OP1 command. In addition if previous OP1 command is Request Monotonic Counter and if SPI Flash returns successful completion extended status then it returns valid values in the Tag, Counter Data and Signature field. Otherwise the values returned in Tag, Counter and Signature field are invalid. The controller may abort the read prematurely prior to completely reading the entire payload. This may occur when the controller wants to simply read the extended status or when it observes an error being returned in the extended status field. The controller may also continue reading past the defined payload size of 49 bytes. Since this is an error condition, the SPI Flash may return any data past the defined payload size. The controller must ignore the data.

Figure 36. Read Data Sequence Diagram





If previous OP1 command is Request Monotonic Counter and if SPI Flash returns successful completion extended status then it returns valid values in the Tag, Counter Data and Signature field. It calculates HMAC-SHA-256 signatures based on following parameters.

- HMAC Message[127:0] = Tag [95:0], Counter_Data_Read[31:0]
- HMAC Key[255:0] = HMAC_Key_Register[Counter_Address][255:0]

Table 17. Extended Status Register Definition

Extended Status[7:0]	Applicable CmdType(s)	Description
00000000	-	Power On State (OP2 issued directly after power-up).
10000000	00, 01, 02, 03	This status is set on successful completion (no errors) of OP1 command.
0xxxxx1	00, 01, 02, 03	This bit is set to 1, when device is busy executing OP1 command. It is reset to 0 when OP1 command execution is done.
0xxxx1x	00, 01	This bit is set only when the correct payload size is received. When cmdtype = 0, this error bit must be set on Root Key Register Overwrite or Counter Address out of range or Truncated Signature mis-match error. For remaining cmdtype = 1 this bit is set when the corresponding monotonic counter is uninitialized
0xxx1xx	01, 02, 03	This bit is set on Signature Mismatch, Counter Address out of range when correct payload size is received; or Cmdtype is out of range; or incorrect payload size is received.
0xxx1xxx	02, 03	This bit is set on HMAC Key Register or Monotonic Counter uninitialized on previous OP1 command when correct payload size is received
0xx1xxxx	02	This bit is set on Counter Data Mismatch on previous increment when correct payload size is received
0x1xxxxx	-	Fatal Error. It is set when no valid counter is found after initialization.
Current value	-	Extended status register will not be updated until first 8 bits of OP1 is received. The correct error type is reflected for any OP1 operation that exceeds a minimum of 16 clocks with active chip-select.

8.6 Operations Allowed/Disallowed During RPMC Operation

In the deep power down state OP1, OP2 commands are ignored until the part comes out of deep power down state.

WREN state does not affect the OP1 command execution inside the SPI Flash.

Suspend operation can be used to execute high-priority reads from the flash device while a long-latency operation is underway. However, OP1 is not recommended when the flash device is in WIP or suspended state.

In the table below, OP1 state is defined as the time starting with a transaction with OP1 op-code sent to the device and ending when the device clears the extended status busy bit. During OP1 state if a suspend transaction is received, the SPI Flash will ignore the suspend command and continue with the execution of the current OP1 command as described in the table below.

P/E state is defined as the time starting with a transaction with write or erase op-code sent to the device and ending when the device clears



the status busy bit. P/E Suspended State starts when the device sets the program suspend status done bit after receiving a program suspend op-code. During P/E State and P/E Suspended State, OP1 is also allowed but not recommended

The table below shows all operation support in each state.

Table 18. RPMC Operation

Operation	OP1 state	P/E state	P/E Suspended State
Suspend	Ignored	Yes-> P/E Suspended State(not chip erase or write status operation) No ->remain P/E state (chip erase or write status operation)	No
Resume	Ignored	No	Yes -> P/E state
All reads except Read status	Yes	No	Yes
All writes/erases	Yes	No	No
OP1	No	Yes but not recommended	Yes but not recommended
Write status	Yes	No	No
OP2	Yes->OP1 busy state (when extended status busy is 1) ->OP1 done state (when extended status busy is 0)	Yes. Will indicate the status associated with the OP1 operation.	Yes. Will indicate the status associated with the OP1 operation
Read status	Yes. Will indicate the busy state associated with the subsequent transaction issued to the SPI Flash.	Yes. Will indicate the busy state associated with the subsequent transaction issued to the SPI Flash.	Yes. Will indicate the busy state associated with the subsequent transaction issued to the SPI Flash.



9 ELECTRICAL CHARACTERISTICS

9.1 Power-On Timing

Figure 37. Power-On Timing Sequence Diagram

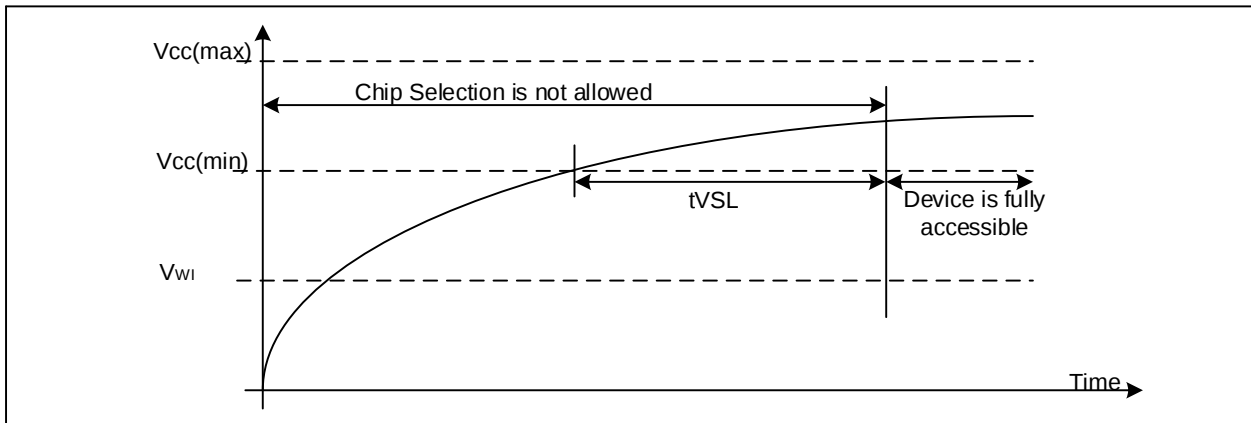


Table 19. Power-Up Timing and Write Inhibit Threshold

Symbol	Parameter	Min.	Max.	Unit
tVSL	VCC (min.) to device operation	1.8		ms
VWI	Write Inhibit Voltage	1.5	2.5	V

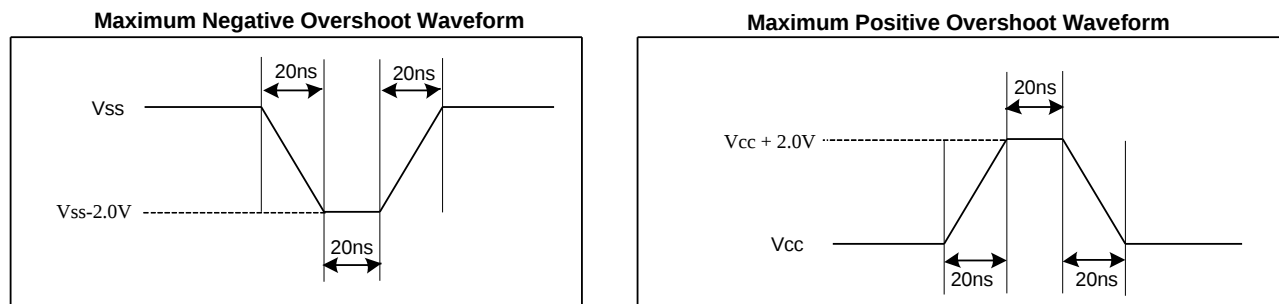
9.2 Initial Delivery State

The device is delivered with the memory array erased: all bits are set to 1 (each Byte contains FFH). The Status Register contains 00H, except that DRV0 bit (S21) and QE bit (S9) are set to 1.

9.3 Absolute Maximum Ratings

Parameter	Value	Unit
Ambient Operating Temperature (T_A)	-40 to 85	°C
Storage Temperature	-65 to 150	°C
Transient Input/Output Voltage (note: overshoot)	-2.0 to VCC+2.0	V
Applied Input/Output Voltage	-0.6 to VCC+0.4	V
VCC	-0.6 to 4.2	V

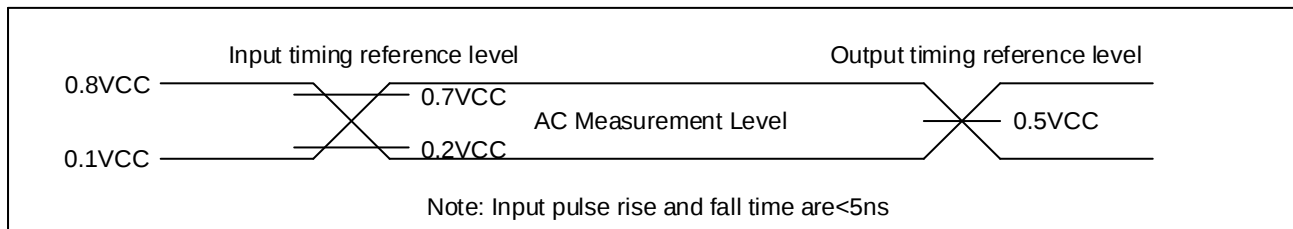
Figure 38. Input Test Waveform and Measurement Level



9.4 Capacitance Measurement Conditions

Symbol	Parameter	Min.	Typ.	Max.	Unit	Conditions
C _{IN}	Input Capacitance			12	pF	V _{IN} =0V
C _{OUT}	Output Capacitance			16	pF	V _{OUT} =0V
C _L	Load Capacitance	30			pF	
	Input Rise And Fall time			5	ns	
	Input Pulse Voltage	0.1VCC to 0.8VCC			V	
	Input Timing Reference Voltage	0.2VCC to 0.7VCC			V	
	Output Timing Reference Voltage	0.5VCC			V	

Figure 39. Absolute Maximum Ratings Diagram





9.5 DC Characteristics

($T_A = -40^{\circ}\text{C} \sim 85^{\circ}\text{C}$, $V_{CC} = 2.7 \sim 3.6\text{V}$)

Symbol	Parameter	Test Condition	Min.	Typ.	Max.	Unit.
I_{LI}	Input Leakage Current				± 4	μA
I_{LO}	Output Leakage Current				± 4	μA
I_{CC1}	Standby Current	$CS\# = V_{CC}$, $VIN = V_{CC}$ or V_{SS}		15	60	μA
I_{CC2}	Deep Power-Down Current	$CS\# = V_{CC}$, $VIN = V_{CC}$ or V_{SS}		2	18	μA
I_{CC3}	Operating Current (Read)	$CLK = 0.1V_{CC} / 0.9V_{CC}$ at 104MHz, $Q = \text{Open}(x4 \text{ I/O})$		12	20	mA
		$CLK = 0.1V_{CC} / 0.9V_{CC}$ at 80MHz, $Q = \text{Open}(x4 \text{ I/O})$		8	13	mA
I_{CC4}	Operating Current (PP)	$CS\# = V_{CC}$		15	25	mA
I_{CC5}	Operating Current (WRSR)	$CS\# = V_{CC}$		15	25	mA
I_{CC6}	Operating Current (SE)	$CS\# = V_{CC}$		15	25	mA
I_{CC7}	Operating Current (BE)	$CS\# = V_{CC}$		15	25	mA
I_{CC8}	Operating Current (CE)	$CS\# = V_{CC}$		15	25	mA
V_{IL}	Input Low Voltage				$0.2V_{CC}$	V
V_{IH}	Input High Voltage		$0.7V_{CC}$			V
V_{OL}	Output Low Voltage	$I_{OL} = 100\mu\text{A}$			0.2	V
V_{OH}	Output High Voltage	$I_{OH} = -100\mu\text{A}$	$V_{CC} - 0.2$			V

Note:

1. Typical value at $T_A = 25^{\circ}\text{C}$, $V_{CC} = 3.3\text{V}$.
2. Value guaranteed by design and/or characterization, not 100% tested in production.



9.6 AC Characteristics

(T_A = -40℃~85℃, VCC=2.7~3.6V)

Symbol	Parameter	Min.	Typ.	Max.	Unit.
f _{CL}	Serial Clock Frequency For: all commands except Read (03H)			104	MHz
f _R	Serial Clock Frequency For: Read (03H)			80	MHz
t _{CLH}	Serial Clock High Time	45% (1/Fc)			ns
t _{CLL}	Serial Clock Low Time	45% (1/Fc)			ns
t _{CLCH}	Serial Clock Rise Time (Slew Rate)	0.2			V/ns
t _{CHCL}	Serial Clock Fall Time (Slew Rate)	0.2			V/ns
t _{SLCH}	CS# Active Setup Time	5			ns
t _{CHSH}	CS# Active Hold Time	5			ns
t _{SHCH}	CS# Not Active Setup Time	5			ns
t _{CHSL}	CS# Not Active Hold Time	5			ns
t _{SHSL}	CS# High Time (Read/Write)	20			ns
t _{SHQZ}	Output Disable Time			6	ns
t _{CLQX}	Output Hold Time	1.2			ns
t _{DVCH}	Data In Setup Time	2			ns
t _{CHDX}	Data In Hold Time	2			ns
t _{CLQV}	Clock Low To Output Valid			7	ns
t _{DP}	CS# High To Deep Power-Down Mode			3	μs
t _{RES1}	CS# High To Standby Mode Without Electronic Signature Read			20	μs
t _{RES2}	CS# High To Standby Mode With Electronic Signature Read			20	μs
t _{SUS}	CS# High To Next Command After Suspend			20	μs
t _{RS}	Latency Between Resume And Next Suspend	100			μs
t _{RST}	CS# High To Next Command After Reset (Except From Erase)			30	μs
t _{RST_E}	CS# High To Next Command After Reset (From Erase)			12	ms
t _W	Write Status Register Cycle Time		5	30	ms
t _{BP1}	Byte Program Time (First Byte)		40	70	μs
t _{BP2}	Additional Byte Program Time (After First Byte)		2.5	12	μs
t _{PP}	Page Programming Time		0.5	2.4	ms
t _{SE}	Sector Erase Time		45	300	ms
t _{BE1}	Block Erase Time (32K Bytes)		0.15	1.2	s
t _{BE2}	Block Erase Time (64K Bytes)		0.25	1.6	s
t _{CE}	Chip Erase Time (GD25R128E)		50	100	s
t _{WRKR}	Write Root key register		3	5.5	ms



t_{UHKR}	Update HMAC Key Register		120		μs
t_{IMC}	Increment Monotonic Counter		20	300	ms
t_{RMC}	Request Monotonic Counter		100	1200	μs

Note:

1. Typical value at $T_A = 25^\circ C$.
2. Value guaranteed by design and/or characterization, not 100% tested in production.

Figure 40. Input Timing

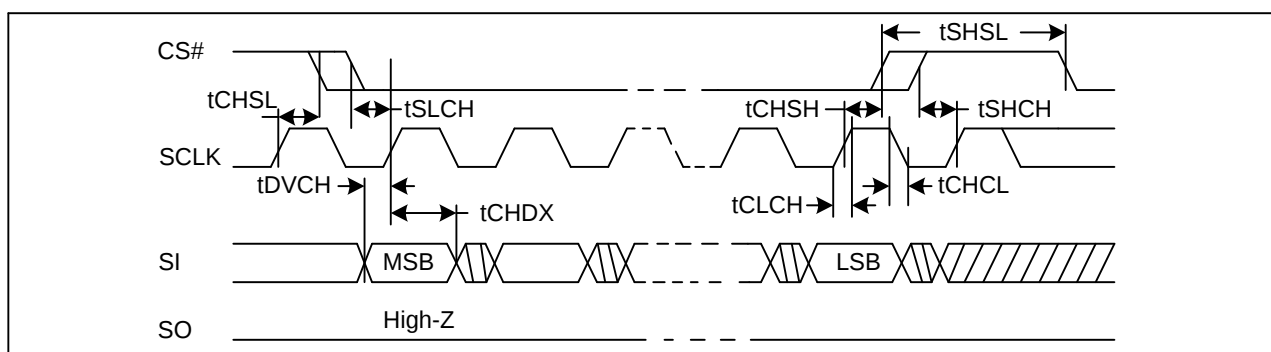


Figure 41. Output Timing

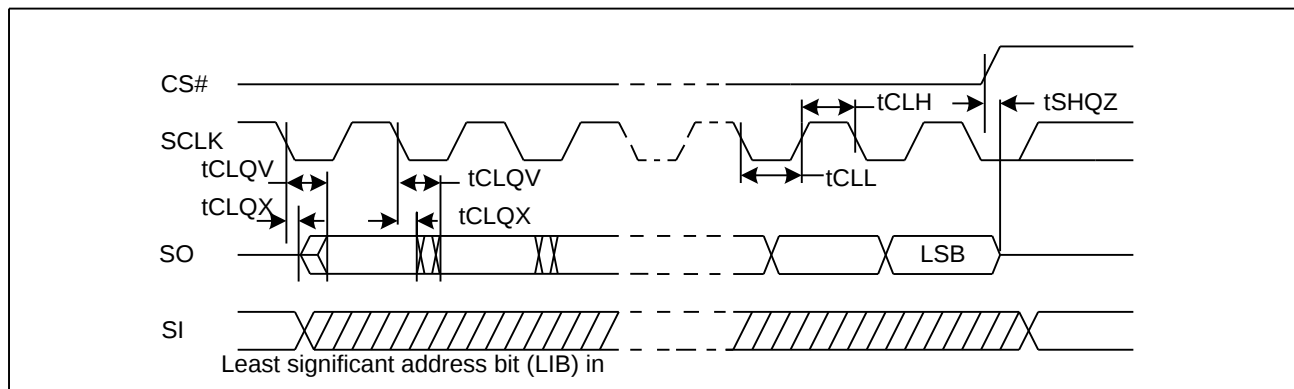
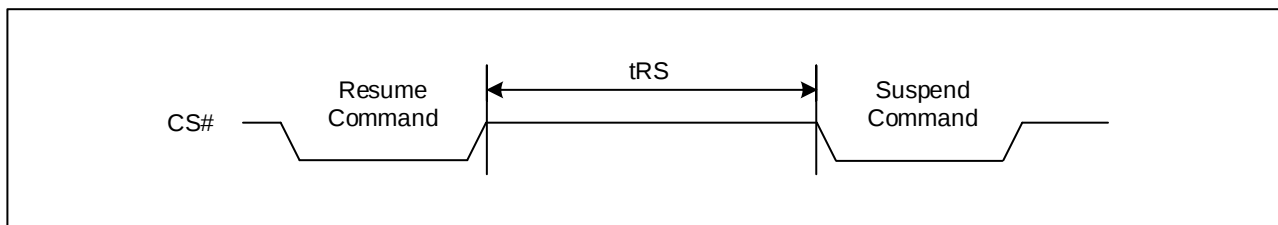
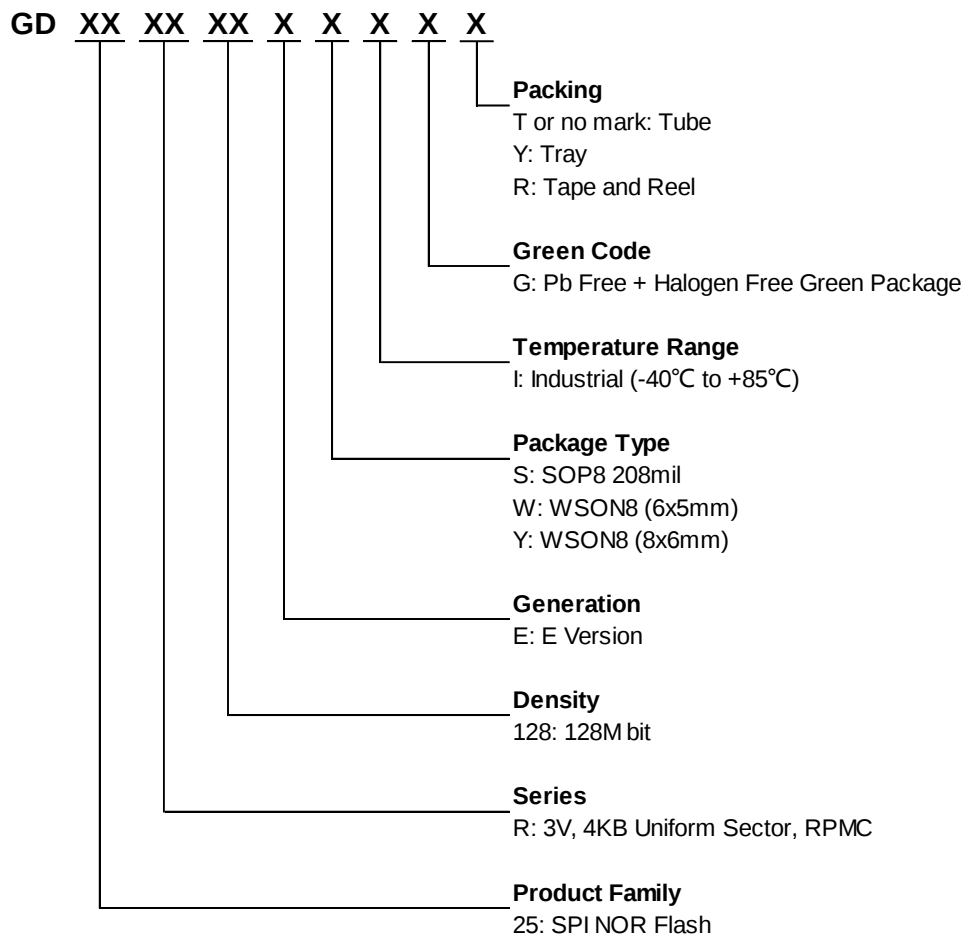


Figure 42. Resume to Suspend Timing Diagram





10 ORDERING INFORMATION





10.1 Valid Part Numbers

Please contact GigaDevice regional sales for the latest product selection and available form factors.

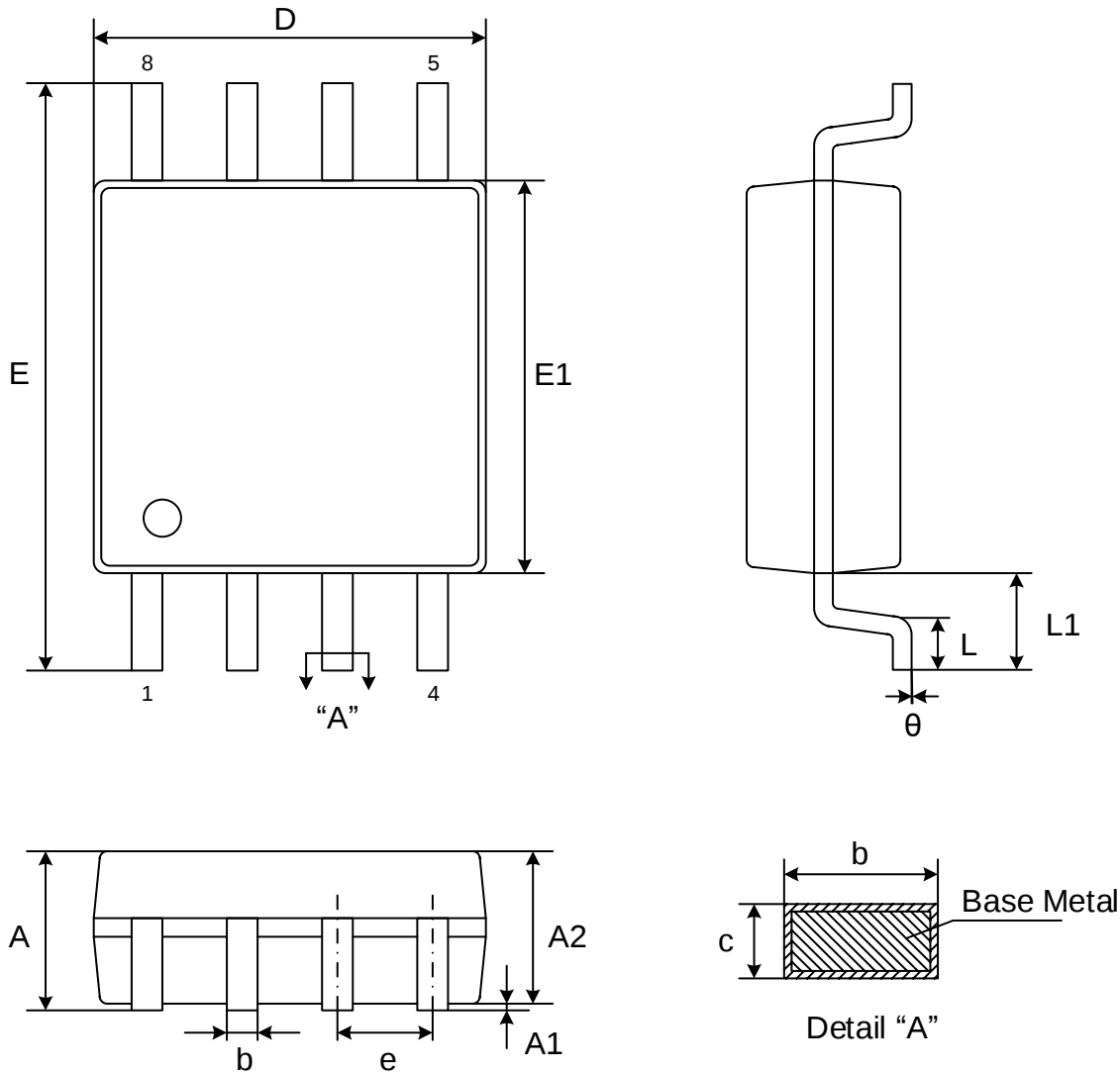
Temperature Range I: Industrial (-40°C to +85°C)

Product Number	Density	Package Type	Packing Options
GD25R128ESIG	128Mbit	SOP8 208mil	T/Y/R
GD25R128EWIG	128Mbit	WSO8 (6x5mm)	Y/R
GD25R128EYIG	128Mbit	WSO8 (8x6mm)	Y/R



11 PACKAGE INFORMATION

11.1 Package SOP8 208MIL



Dimensions

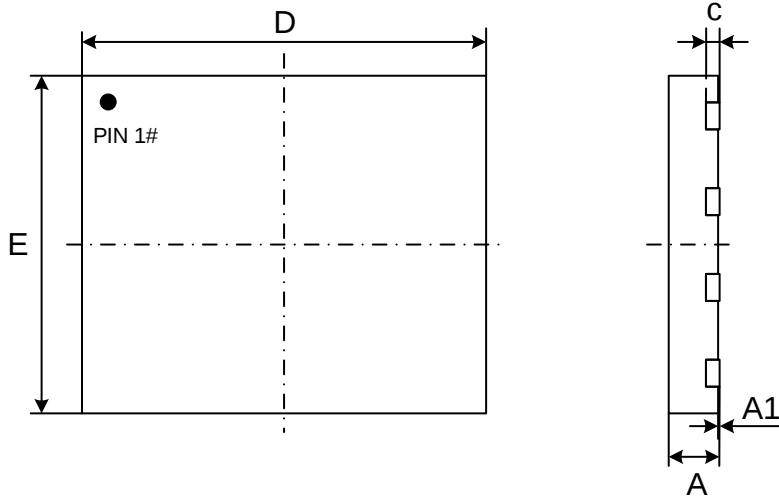
Symbol		A	A1	A2	b	c	D	E	E1	e	L	L1	θ
Unit													
mm	Min	-	0.05	1.70	0.31	0.15	5.13	7.70	5.18	1.27	0.50	1.31	0°
	Nom	-	0.15	1.80	0.41	0.20	5.23	7.90	5.28		-		-
	Max	2.16	0.25	1.90	0.51	0.25	5.33	8.10	5.38		0.85		8°

Note:

- Both the package length and width do not include the mold flash.
- Seating plane: Max. 0.1mm.

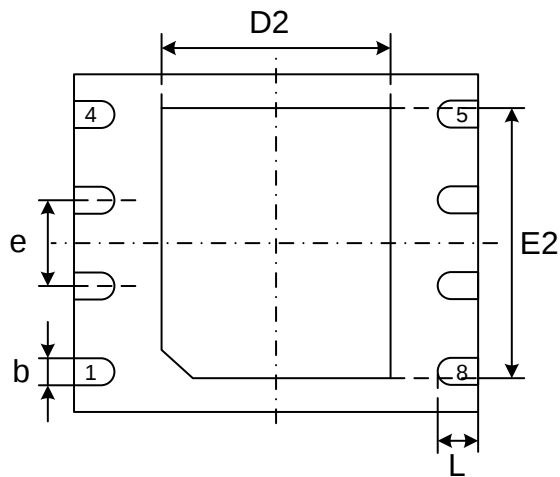


11.2 Package WSON8 (6x5mm)



Top View

Side View



Bottom View

Dimensions

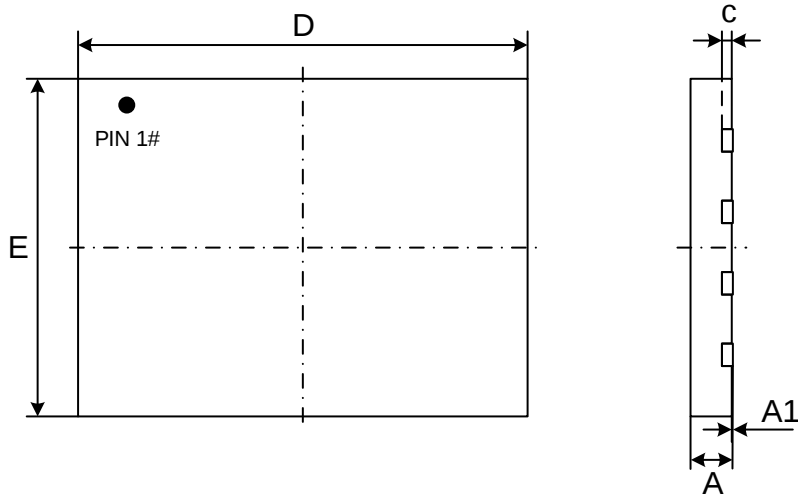
Symbol		A	A1	c	b	D	D2	E	E2	e	L
Unit											
mm	Min	0.70	0.00	0.180	0.35	5.90	3.30	4.90	3.90	1.27	0.50
	Nom	0.75	0.02	0.203	0.40	6.00	3.40	5.00	4.00		0.60
	Max	0.80	0.05	0.250	0.50	6.10	3.50	5.10	4.10		0.75

Note:

- Both the package length and width do not include the mold flash.
- The exposed metal pad area on the bottom of the package is floating.
- Coplanarity $\leq 0.08\text{mm}$. Package edge tolerance $\leq 0.10\text{mm}$.
- The lead shape may be of little difference according to different package lead frames. These lead shapes are compatible with each other.

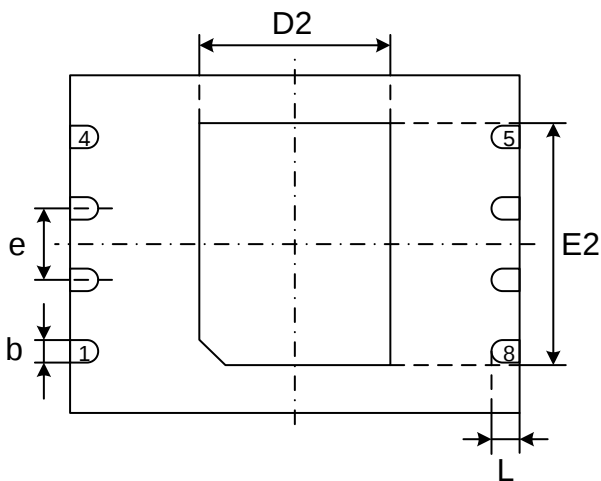


11.3 Package WSON8 (8x6mm)



Top View

Side View



Bottom View

Dimensions

Symbol		A	A1	c	b	D	D2	E	E2	e	L
Unit											
mm	Min	0.70	0.00	0.180	0.35	7.90	3.30	5.90	4.20	1.27	0.45
	Nom	0.75	0.02	0.203	0.40	8.00	3.40	6.00	4.30		0.50
	Max	0.80	0.05	0.250	0.45	8.10	3.50	6.10	4.40		0.55

Note:

- Both the package length and width do not include the mold flash.
- The exposed metal pad area on the bottom of the package is floating.
- Coplanarity $\leq 0.08\text{mm}$. Package edge tolerance $\leq 0.10\text{mm}$.
- The lead shape may be of little difference according to different package lead frames. These lead shapes are compatible with each other.



12 REVISION HISTORY

Version No	Description	Page	Date
1.0	Initial release	All	2020-10-21



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